## **Use cases**

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#### CHAPTER

## ONE

## **USE CASES**

Here are several toulbar2 use cases, where toulbar2 has been used in order to resolve different problems. According to cases, they can be used to overview, learn, use toulbar2... They may contain source code, explanations, possibility to run yourself...

You will find the mentioned examples, among the following exhaustive list of examples.

#### toulbar2 and Deep Learning :

- Visual Sudoku Tutorial
- Visual Sudoku Application

Some applications based on toulbar2 :

- Mendelsoft : Mendelsoft detects Mendelian errors in complex pedigree [Sanchez et al, Constraints 2008].
- **Pompd** : POsitive Multistate Protein Design, [Vucini et al Bioinformatics 2020]
- Visual Sudoku Application

Misc :

• A sudoku code

#### CHAPTER

TWO

## LIST OF ALL EXAMPLES

# 2.1 Sudoku puzzle in Pytoulbar2

The Sudoku is a widely known puzzle game that consists in filling a grid with numbers. The typical grid has 9x9 cells in total and each cell must contain an integer between 1 and 9.

Additional constraints need to be enforced to solve the puzzle. In each line, the same number cannot appear twice. The same type of constraint occurs for each column. Finally, each sub square of size 3x3 located every 3 cells must also not contain duplicates. The image bellow shows an example of a Sudoku grid given with its initial values. The goal is to deduce the other values while verifying the different constraints.



## 2.1.1 Getting started

Before starting, make sure pytoulbar2 is installed in your environment. It can be installed via the command:

#### pip install pytoulbar2

We first create a CFN object. CFN, which stands for Cost Function Network, is the main object that is manipulated in pytoulbar2. It represents a problem to solve expressed as a network of cost functions, i.e a set of variables that are connected to each other through discrete cost functions (or constraints).

As ToulBar2 is an optimization framework, an optional upper bound can be provided to the CFN object in order to exclude any solution whose value exceeds this bound. In the case of a Sudoku puzzle, since the problem does not contain a numerical objective, an upper bound of 1 can be chosen.

import pytoulbar2

```
top = 1 # upper bound value of the problem
cfn = pytb2.CFN(top)
```

#### 2.1.2 Representing the grid in ToulBar2

To represent our problem in pytoulbar2, it is necessary to define discrete decision variables. The variables will represent the various choices that can be made to build a solution to the problem. In the Sudoku puzzle, decision variables are typically the different cells of the grid. Their values would be the possible integers they can be assigned to, from 1 to 9. We use the function AddVariable of our cfn object to create the variables.

```
# variable creation
for row in range(9):
    for col in range(9):
        cfn.AddVariable('cell_'+str(row)+'_'+str(col), range(9))
```

#### 2.1.3 Solving first the grid

It is already possible to "solve" the puzzle with ToulBar2, as the cfn object contains the variables of the problem. The function Solve is used to run the solving algorithm.

result = cfn.Solve(showSolutions = 3)

If ToulBar2 finds a solution, it will return an array containing the chosen values for each variable:

print(result[0][0])

Would return 0, meaning the chosen value of the first variable (upper left cell) is 1 (first value of the list). We then define a function to print the solutions as a grid:

```
# print a solution as a Sudoku grid
def print_grid(solution):
  print('-----')
  var_index = 0
  for row in range(9):
     line = ''
     for col in range(9):
        if col % 3 == 0:
             line += '|'
        line += ' '
        line += str(solution[var_index]+1)
        if col % 3 == 2:
             line += ' '
        var_index += 1
     line += '|'
     print(line)
     if row % 3 == 2:
        print('-----')
```

## # print the first solution print\_grid(result)

Which helps to display the first solution:

#### 2.1.4 Adding initial values

The next step consists in initializing the variables that correspond to cells for which the value is known. We will use the values in the grid example above, defined as a double array (where 0 means the value is unspecified) :

```
# define known values
initial_values = [[5,3,0,0,7,0,0,0,0],
        [6,0,0,1,9,5,0,0,0],
        [0,9,8,0,0,0,0,6,0],
        [8,0,0,0,6,0,0,0,3],
        [4,0,0,8,0,3,0,0,1],
        [7,0,0,0,2,0,0,0,6],
        [0,6,0,0,0,0,2,8,0],
        [0,0,0,4,1,9,0,0,5],
        [0,0,0,0,8,0,0,7,9]]
```

Variables can be assigned with the function Assign. The variable and its value can be specified as integer indexes or as strings.

```
var_index = 0
for row in range(9):
    for col in range(9):
        if initial_values[row][col] != 0:
            cfn.Assign(var_index, initial_values[row][col]-1)
        var_index += 1
result = cfn.Solve(showSolutions = 3)
print_grid(result[0])
```

#### **\*** Caution

Although we have already solved the problem once, a CFN object cannot execute its Solve function twice in a row. The object must be recreated or the function must be called only once.

The solution returned by the algorithm this time looks like this :

The initial values are now correctly specified in the variables.

#### 2.1.5 Adding constraints and solving the grid

The last missing part before being able to compute a solution is the constraints. Starting with the row constraints, we must ensure that none of the variables in the same row will be assigned to the same values. This constraint is usually called *all different* and can be added with the function AddAllDifferent. The function takes as an argument a list of indices of the variables that must differ. The constraint on the first row is obtained via:

```
cfn.AddAllDifferent([var_ind for var_ind in range(9)])
```

Which generates the following first row in the solution:

```
| 5 3 1 | 2 7 4 | 9 8 6 |
```

Constraints for each row can be added by varying the column index for each row:

```
# row constraints
for row_ind in range(9):
    cfn.AddAllDifferent([row_ind*9+col_ind for col_ind in range(9)])
```

Constraints for each column are obtained similarly:

```
# column constraints
for col_ind in range(9):
    cfn.AddAllDifferent([row_ind*9+col_ind for row_ind in range(9)])
```

At this point, the solution is not correct yet since sub-grids of size 3x3 may contain duplicates, such as the values 9 and 3 in the example below:

| 5 3 9 |

| 6 2 3 | | 1 9 8 |

Additional constraints are added for each of the 9 sub-grids:

```
# sub grids constraints
for sub_ind1 in range(3): # row offset
    for sub_ind2 in range(3): # column offset
        cfn.AddAllDifferent([(sub_ind1*3+row_ind)*9+ sub_ind2*3+col_ind for col_ind in_
        -range(3) for row_ind in range(3)])
```

These last constraints allow to obtain a consistent solution to the Sudoku puzzle finally:

#### 2.1.6 Conclusion

This short introduction shows how to represent a problem in ToulBar2 via its python interface Pytoulbar2 by defining the problem variables and constraints and how to obtain a solution to this problem. Below is the complete python script from this tutorial.

```
sudoku_tutorial.py
```

```
import pytoulbar2 as pytb2
# print a solution as a sudoku grid
def print_grid(solution):
    print('------')
    var_index = 0
    for row in range(9):
        line = ''
        for col in range(9):
            if col % 3 == 0:
                line += '|'
                line += '|'
                line += str(solution[var_index]+1)
                if col % 3 == 2:
```

```
line += ' '
           var_index += 1
       line += '|'
       print(line)
       if row % 3 == 2:
           print('-----')
cfn = pytb2.CFN()
# variables
for row in range(9):
   for col in range(9):
        cfn.AddVariable('cell_'+str(row)+'_'+str(col), range(9))
# define known values
initial_values = [[5,3,0,0,7,0,0,0,0],
                [6,0,0,1,9,5,0,0,0],
                [0,9,8,0,0,0,0,6,0],
                [8,0,0,0,6,0,0,0,3],
                [4,0,0,8,0,3,0,0,1],
                [7,0,0,0,2,0,0,0,6],
                [0,6,0,0,0,0,2,8,0],
                [0,0,0,4,1,9,0,0,5],
                [0,0,0,0,8,0,0,7,9]]
var_index = 0
for row in range(9):
    for col in range(9):
       if initial_values[row][col] != 0:
            cfn.Assign(var_index, initial_values[row][col]-1)
       var_index += 1
# row constraints
for row_ind in range(9):
    cfn.AddAllDifferent([row_ind*9+col_ind for col_ind in range(9)])
# column constraints
for col_ind in range(9):
   cfn.AddAllDifferent([row_ind*9+col_ind for row_ind in range(9)])
# sub grids constraints
for sub_ind1 in range(3): # row offset
    for sub_ind2 in range(3): # column offset
        cfn.AddAllDifferent([(sub_ind1*3+row_ind)*9+ sub_ind2*3+col_ind for col_ind in_

→range(3) for row_ind in range(3)])
result = cfn.Solve(showSolutions = 3, allSolutions=1)
print_grid(result[0])
```

# 2.2 Sudoku puzzle with libtb2 in C++

The Sudoku is a widely known puzzle game that consists in filling a grid with numbers. The typical grid has 9x9 cells in total and each cell must contain an integer between 1 and 9.

Additional constraints need to be enforced to solve the puzzle. In each line, the same number cannot appear twice. The same type of constraint occurs for each column. Finally, each sub-square of size 3x3 located every 3 cells must also not contain duplicates. The image below shows an example of a Sudoku grid given with its initial values. The goal is to deduce the other values while verifying the different constraints.



#### 2.2.1 Getting started

Before starting, make sure the ToulBar2 C++ library binaries are installed in your system (libtb2.so, see installation section from sources or binaries for more instructions).

We first create a WeightedCSPSolver object. This object is in charge of executing the algorithm to solve the Sudoku grid and internally creates a WeightedCSP object to store the problem.

WeightedCSP objects are used in ToulBar2 to define the optimization or decision problems. The problem is expressed as a set of discrete variables that are connected to each other through cost functions (or constraints).

As ToulBar2 is an optimization framework, an optional upper bound can be provided to the solver object in order to exclude any solution whose value exceeds this bound. In the case of a Sudoku puzzle, since the problem does not contain a numerical objective, an upper bound of 1 can be chosen.

In order to compile the following code, assuming we are in the main ToulBar2 source repository, the same compilation flags as used to compile libtb2.so must be used: g++ -DBOOST -DLONGDOUBLE\_PROB -DLONGLONG\_COST -I./ src -o sudoku sudoku\_tutorial.cpp libtb2.so

```
#include <iostream>
#include <iostream>
using namespace std;
int main() {
    // initialization
    tb2init();
```

```
initCosts();
Cost top = 1;
// creation of the solver object
WeightedCSPSolver* solver = WeightedCSPSolver::makeWeightedCSPSolver(top);
// access to the wcsp object created by the solver
WeightedCSP* wcsp = solver->getWCSP();
delete solver;
return 0;
```

#### 2.2.2 Representing the grid in ToulBar2

To represent our problem in pytoulbar2, it is necessary to define discrete decision variables. The variables will represent the various choices that can be made to build a solution to the problem. In the Sudoku puzzle, decision variables are typically the different cells of the grid. Their values would be the possible integers they can be assigned to, from 1 to 9. We use the makeEnumeratedVariable function of the wcsp object to make the variables.

#### 2.2.3 Solving first the grid

It is already possible to solve the puzzle with ToulBar2, as the cfn object contains the variables of the problem. The WeightedCSPSolver::solve function is used to run the solving algorithm.

```
// close the model definition
wcsp->sortConstraints();
// solve the problem
bool hasSolution = solver->solve();
```

#### 🛕 Warning

}

It is important to systematically call the function WeightedCSP::sortConstraints before solving the problem to close the model definition.

Problems may sometimes be especially hard to solve. In such cases, a timeout can be set on the solver to stop the search before the optimality proof is complete, or before a solution is found at all. When doing so, ToulBar2 throws an exception that must be caught to properly clean the problem data structures:

```
try {
   bool hasSolution = solver->solve();
} catch(const exception& ex) {
   cout << "no solution found: " << ex.what() << endl;
}</pre>
```

When ToulBar2 returns a solution, the solution can be accessed as a std::vector of Value, specifying the value that is assigned to each variable of the problem (in the same order they were defined):

```
if(hasSolution) {
    std::vector<Value> solution = solver->getSolution();
    cout << "the first value is " << solution[0] << endl;
}</pre>
```

The above code would display 0, meaning the chosen value of the first variable (upper left cell) is 1 (first value of the list). We then define a function to print the solution as a grid :

```
// print a solution as a grid
void printSolution(const std::vector<Value>& solution) {
   cout << "-----" << endl;</pre>
   size_t var_index = 0;
  for(size_t row = 0; row < 9; row ++) {</pre>
      for(size_t col = 0; col < 9; col ++) {</pre>
         if(col % 3 == 0) {
               cout << "|";
         }
         cout << " " << to_string(solution[var_index]+1);</pre>
         if(col % 3 == 2) {
               cout << " ";
         }
         var_index += 1;
         }
         cout << "|" << endl;</pre>
         if(row % 3 == 2) {
            cout << "----
                           -----" << endl;
         }
      }
}
```

This function helps to visualize the variables' values as a real Sudoku grid, as follows :

```
// print the first solution
if(hasSolution) {
    std::vector<Value> solution = solver->getSolution();
```

```
printSolution(solution)
```

#### }

Which helps to display the first solution:

#### 2.2.4 Assignment to the input values

The next step consists in initializing the variables that correspond to cells for which the value is known. We will use the values in the grid example above, defined as a two-dimensional vector (where 0 means the value is unspecified):

// grid data									
<pre>std::vector<std::vector<value> &gt; input_grid {</std::vector<value></pre>	<b>{5</b> ,	3,	0,	0,	7,	0,	0,	0,	≬},
	{ <mark>6</mark> ,	0,	0,	1,	9,	5,	0,	0,	≬},
	{≬,	9,	8,	0,	0,	0,	0,	6,	≬},
	<b>{</b> 8,	0,	0,	0,	6,	0,	0,	0,	3},
	{4,	0,	0,	8,	0,	3,	0,	0,	1},
	{7,	0,	0,	0,	2,	0,	0,	0,	6},
	{≬,	6,	0,	0,	0,	0,	2,	8,	≬},
	{≬,	0,	0,	4,	1,	9,	0,	0,	5},
	{ <b>0</b> ,	0,	0,	0,	8,	0,	0,	7,	9} };

Variables can be assigned with the WeightedCSP::assign function. The variable and its value can be specified as integer indexes or as strings.

```
// input values initialization
size_t var_ind = 0;
for(size_t row = 0; row < 9; row ++) {
    for(size_t col = 0; col < 9; col ++) {
        if(input_grid[row][col] != 0) {
            wcsp->assign(var_ind, input_grid[row][col]-1);
        }
        var_ind ++;
    }
}
```

#### **A** Warning

Although we already solved the problem once, a WeightedCSP object cannot execute its WeightedCSPSolver::solve function twice in a row. The object must be recreated or the function must be called only once.

The solution returned by the algorithm this time looks like this:

The initial values are now correctly specified in the variables.

#### 2.2.5 Adding constraints and solving the grid

The missing part to be able to generate a solution is the constraints. Starting with the row constraints, we must ensure that none of the variables in the same row will be assigned to the same values. This constraint is usually called *all different* and can be added with the WeightedCSPSolver::postWAllDiff function. The function takes as arguments a list of indices of the variables that must differ (the scope of the constraint) as well as two parameters specifying how the constraint is encoded, which we do not further describe in this tutorial. We start by adding a constraint for the first row:

```
std::vector<int> scope = {0,1,2,3,4,5,6,7,8};
std::string semantics = "hard";
std::string prop = "knapsack";
wcsp->postWAllDiff(scope, semantics, prop, top);
```

Which generates the following first row in the solution :

```
| 5 3 1 | 2 7 4 | 6 8 9 |
```

Constraints for each row can be added by varying the column index for each row :

```
// add one "all different" constraint for each row
for(int row = 0; row < 9; row ++) {
   std::vector<int> row_scope;
   for(int col = 0; col < 9; col ++) {
      row_scope.emplace_back(row*9+col);
   }</pre>
```

```
wcsp->postWAllDiff(row_scope, semantics, prop, top);
```

Constraints for each column are obtained similarly:

```
// add one "all different" constraint for each column
for(int col = 0; col < 9; col ++) {
   std::vector<int> col_scope;
   for(int row = 0; row < 9; row ++) {
      col_scope.emplace_back(row*9+col);
   }
   wcsp->postWAllDiff(col_scope, semantics, prop, top);
}
```

At this point, the solution is not correct yet since sub-grids of size 3x3 may contain duplicates, such as the values 9 and 3 in the example below:

| 9 7 6 | | 1 9 5 | | 5 4 2 |

}

A set of 9 additional allDifferent constraints can be defined to finalize our Sudoku model definition:

```
// add one "all different" constraint for each 3x3 sub-grid
for(int sub_ind1 = 0; sub_ind1 < 3; sub_ind1 ++) {
   for(int sub_ind2 = 0; sub_ind2 < 3; sub_ind2 ++) {
      std::vector<int> sub_scope;
      for(int row_ind = 0; row_ind < 3; row_ind ++) { // iterate inside the 3x3 sub-grid
      for(int col_ind = 0; col_ind < 3; col_ind ++) {
         sub_scope.emplace_back((sub_ind1*3+row_ind)*9+sub_ind2*3+col_ind);
      }
    }
    wcsp->postWAllDiff(sub_scope, semantics, prop, top);
    }
}
```

These last constraints allow to finally obtain a consistent solution to the Sudoku puzzle :

#### 2.2.6 Conclusion

This short introduction shows how to represent a problem in ToulBar2 via its C++ library libtb2 by defining the problem variables and constraints and how to obtain a solution to this problem. Below is the complete C++ source code from this tutorial.

```
sudoku_tutorial.cpp
```

```
#include <iostream>
#include <toulbar2lib.hpp>
using namespace std;
// print a solution as a grid
void printSolution(const std::vector<Value>& solution) {
  cout << "-----" << endl;
  size_t var_index = 0;
  for(size_t row = 0; row < 9; row ++) {</pre>
     for(size_t col = 0; col < 9; col ++) {</pre>
        if(col % 3 == 0) {
           cout << "|";
        }
        cout << " " << to_string(solution[var_index]+1);</pre>
        if(col % 3 == 2) {
           cout << " ";
        }
        var_index += 1;
     }
     cout << "|" << endl;</pre>
     if(row % 3 == 2) {
        cout << "-----" << endl;
     }
  }
}
int main() {
   // grid data
  \{6, 0, 0, 1, 9, 5, 0, 0, 0\},\
                                               \{0, 9, 8, 0, 0, 0, 0, 6, 0\},\
                                               \{8, 0, 0, 0, 6, 0, 0, 0, 3\},\
                                               \{4, 0, 0, 8, 0, 3, 0, 0, 1\},\
                                               \{7, 0, 0, 0, 2, 0, 0, 0, 6\},\
                                               \{0, 6, 0, 0, 0, 0, 2, 8, 0\},\
                                               \{0, 0, 0, 4, 1, 9, 0, 0, 5\},\
                                               \{0, 0, 0, 0, 8, 0, 0, 7, 9\}\};
```

```
// initialisation
  tb2init();
  initCosts();
  Cost top = 1;
  // creation of the solver object
  WeightedCSPSolver* solver = WeightedCSPSolver::makeWeightedCSPSolver(top);
  // access to the wcsp object created by the solver
  WeightedCSP* wcsp = solver->getWCSP();
  //-----
  // problem variables
  //-----
  // variable creation
  for(size_t row = 0; row < 9; row ++) {</pre>
     for(size_t col = 0; col < 9; col ++) {</pre>
        wcsp->makeEnumeratedVariable("Cell_" + to_string(row) + "," + to_string(col), 0,
↔ 8);
    }
  }
  // input values initialisation
  size_t var_ind = 0;
  for(size_t row = 0; row < 9; row ++) {</pre>
     for(size_t col = 0; col < 9; col ++) {</pre>
       if(input_grid[row][col] != 0) {
          wcsp->assign(var_ind, input_grid[row][col]-1);
        }
       var_ind ++;
     }
  }
  //-----
  // constraints definition
  //-----
  // all different constraint parameters
  std::string semantics = "hard";
  std::string prop = "knapsack";
  // add one "all different" constraint for each row
  for(int row = 0; row < 9; row ++) {
     std::vector<int> row_scope;
     for(int col = 0; col < 9; col ++) {
        row_scope.emplace_back(row*9+col);
     }
```

```
wcsp->postWAllDiff(row_scope, semantics, prop, top);
  }
  // add one "all different" constraint for each column
   for(int col = 0; col < 9; col ++) {
      std::vector<int> col_scope;
      for(int row = 0; row < 9; row ++) {</pre>
        col_scope.emplace_back(row*9+col);
     }
     wcsp->postWAllDiff(col_scope, semantics, prop, top);
  }
  // add one "all different" constraint for each 3x3 sub square
   for(int sub_ind1 = 0; sub_ind1 < 3; sub_ind1 ++) {</pre>
      for(int sub_ind2 = 0; sub_ind2 < 3; sub_ind2 ++) {</pre>
         std::vector<int> sub_scope;
         for(int row_ind = 0; row_ind < 3; row_ind ++) { // iterate inside the 3x3 sub_</pre>
\rightarrow grid
            for(int col_ind = 0; col_ind < 3; col_ind ++) {</pre>
               sub_scope.emplace_back((sub_ind1*3+row_ind)*9+sub_ind2*3+col_ind);
            }
        }
        wcsp->postWAllDiff(sub_scope, semantics, prop, top);
     }
  }
   //-----
   // solution
   //-----
                        _____
  // close the model definition
  wcsp->sortConstraints();
  // solve the problem
  bool hasSolution = solver->solve();
  if(hasSolution) {
      std::vector<Value> solution = solver->getSolution();
     printSolution(solution);
  }
  // clean the solver object
  delete solver;
  return 0;
}
```

## 2.3 Weighted n-queen problem

## 2.3.1 Brief description

The problem consists in assigning N queens on a NxN chessboard with random costs in (1..N) associated to every cell such that each queen does not attack another queen and the sum of the costs of queen's selected cells is minimized.

## 2.3.2 CFN model

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A solution must have only one queen per column and per row. We create N variables for every column with domain size N to represent the selected row for each queen. A clique of binary constraints is used to express that two queens cannot be on the same row. Forbidden assignments have cost  $k=N^{**}2+1$ . Two other cliques of binary constraints are used to express that two queens do not attack each other on a lower/upper diagonal. We add N unary cost functions to create the objective function with random costs on every cell.

## 2.3.3 Example for N=4 in JSON .cfn format

<section-header><section-header><section-header><section-header><section-header><section-header>

```
problem: { "name": "4-queen", "mustbe": "<17" },</pre>
    variables: {"Q0":["Row0", "Row1", "Row2", "Row3"], "Q1":["Row0", "Row1", "Row2", "Row3"], "Q1":["Row0", "Row1", "Row2", "Row3"], "Q1":["Row0", "Row1", "Row2", "Row3"], "Q1":["Row0", "Row1", "Row1", "Row3"], 
→"],
                                    "Q2":["Row0", "Row1", "Row2", "Row3"], "Q3":["Row0", "Row1", "Row2", "Row3
\rightarrow"]},
    functions: {
         {scope: ["Q0", "Q1"], "costs": [17, 0, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17]},
         {scope: ["Q0", "Q2"], "costs": [17, 0, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17]},
                                              "Q3"], "costs": [17, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17]},
         {scope: ["Q0",
         {scope: ["Q1", "Q2"], "costs": [17, 0, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17]},
         {scope: ["Q1", "Q3"], "costs": [17, 0, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17]},
         {scope: ["Q2", "Q3"], "costs": [17, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17]},
         {scope: ["Q0", "Q1"], "costs": [0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0]},
         {scope: ["Q0", "Q2"], "costs": [0, 0, 0, 0, 0, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0, 0]},
         {scope: ["Q0", "Q3"], "costs": [0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 17, 0, 0, 0]},
         {scope: ["Q1", "Q2"], "costs": [0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0]},
         {scope: ["Q1", "Q3"], "costs": [0, 0, 0, 0, 0, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0, 0]},
         {scope: ["Q2", "Q3"], "costs": [0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0, 0, 0, 0, 17, 0]},
```

{sc	cope:	["Q0",	"Q1"],	"costs":	[0,	17, 0, 0,	0,	0,	17, 0, 0, 0, 0, 17, 0, 0, 0, 0]
{sc	cope:	["Q0",	"Q2"],	"costs":	[0,	0, 17, 0,	0,	0,	0, 17, 0, 0, 0, 0, 0, 0, 0, 0]},
{sc	cope:	["Q0",	"Q3"],	"costs":	[0,	0, 0, 17,	0,	0,	0, 0, 0, 0, 0, 0, 0, 0, 0, 0]},
{sc	cope:	["Q1",	"Q2"],	"costs":	[0,	17, 0, 0,	0,	0,	17, 0, 0, 0, 0, 17, 0, 0, 0, 0]
{sc	cope:	["Q1",	"Q3"],	"costs":	[0,	0, 17, 0,	0,	0,	0, 17, 0, 0, 0, 0, 0, 0, 0, 0]},
{sc	cope:	["Q2",	"Q3"],	"costs":	[0,	17, 0, 0,	0,	0,	17, 0, 0, 0, 0, 17, 0, 0, 0, 0]
{sc	cope:	["Q0"]	, "cost	s": [4, 4	, 3,	4]},			
{sc	cope:	["Q1"]	, "cost	s": [4, 3	, 4,	4]},			
{sc	cope:	["Q2"]	, "cost:	s": [2, 1	, 3,	2]},			
{sc	cope:	["Q3"]	, "cost	s": [1, 2	, 3,	4]}}			
}									

Optimal solution with cost 11 for the 4-queen example :



#### 2.3.4 Python model

The following code using the pytoulbar2 library solves the weighted N-queen problem with the first argument being the number of queens N (e.g. "python3 weightedqueens.py 8").

weightedqueens.py

```
import sys
from random import seed, randint
seed(123456789)
import pytoulbar2
N = int(sys.argv[1])
top = N**2 +1
Problem = pytoulbar2.CFN(top)
for i in range(N):
        Problem.AddVariable('Q' + str(i+1), ['row' + str(a+1) for a in range(N)])
```

```
for i in range(N):
        for j in range(i+1,N):
                #Two queens cannot be on the same row constraints
                ListConstraintsRow = []
                for a in range(N):
                        for b in range(N):
                                if a != b :
                                        ListConstraintsRow.append(0)
                                else:
                                         ListConstraintsRow.append(top)
                Problem AddFunction([i, j], ListConstraintsRow)
                #Two queens cannot be on the same upper diagonal constraints
                ListConstraintsUpperD = []
                for a in range(N):
                        for b in range(N):
                                if a + i != b + j :
                                        ListConstraintsUpperD.append(0)
                                else:
                                         ListConstraintsUpperD.append(top)
                Problem AddFunction([i, j], ListConstraintsUpperD)
                #Two queens cannot be on the same lower diagonal constraints
                ListConstraintsLowerD = []
                for a in range(N):
                        for b in range(N):
                                if a - i != b - j :
                                        ListConstraintsLowerD.append(0)
                                else:
                                         ListConstraintsLowerD.append(top)
                Problem AddFunction([i, j], ListConstraintsLowerD)
#Random unary costs
for i in range(N):
        ListConstraintsUnaryC = []
        for j in range(N):
                ListConstraintsUnaryC.append(randint(1,N))
        Problem AddFunction([i], ListConstraintsUnaryC)
#Problem.Dump('WeightQueen.cfn')
Problem.CFN.timer(300)
res = Problem.Solve(showSolutions = 3)
if res:
        for i in range(N):
                row = ['X' if res[0][j]==i else ' ' for j in range(N)]
                print(row)
        # and its cost
        print("Cost:", int(res[1]))
```

## 2.4 Weighted latin square problem

## 2.4.1 Brief description

The problem consists in assigning a value from 0 to N-1 to every cell of a NxN chessboard. Each row and each column must be a permutation of N values. For each cell, a random cost in (1...N) is associated to every domain value. The objective is to find a complete assignment where the sum of the costs associated to the selected values for the cells is minimized.

### 2.4.2 CFN model

We create NxN variables, one for every cell, with domain size N. An AllDifferent hard global constraint is used to model a permutation for every row and every column. Its encoding uses knapsack constraints. Unary cost functions containing random costs associated to domain values are generated for every cell. The worst possible solution is when every cell is associated with a cost of N, so the maximum cost of a solution is N\*\*3, so forbidden assignments have cost  $k=N^{*}3+1$ .

#### 2.4.3 Example for N=4 in JSON .cfn format

```
{
  problem: { "name": "LatinSquare4", "mustbe": "<65" },</pre>
  variables: {"X0_0": 4, "X0_1": 4, "X0_2": 4, "X0_3": 4, "X1_0": 4, "X1_1": 4, "X1_2":
→4, "X1_3": 4, "X2_0": 4, "X2_1": 4, "X2_2": 4, "X2_3": 4, "X3_0": 4, "X3_1": 4, "X3_2
\rightarrow": 4, "X3_3": 4},
  functions: {
    {scope: ["X0_0", "X0_1", "X0_2", "X0_3"], "type:" salldiff, "params": {"metric": "var
\rightarrow", "cost": 65}},
   {scope: ["X1_0", "X1_1", "X1_2", "X1_3"], "type:" salldiff, "params": {"metric": "var
\rightarrow", "cost": 65}},
    {scope: ["X2_0", "X2_1", "X2_2", "X2_3"], "type:" salldiff, "params": {"metric": "var
\rightarrow", "cost": 65}},
    {scope: ["X3_0", "X3_1", "X3_2", "X3_3"], "type:" salldiff, "params": {"metric": "var
\rightarrow", "cost": 65}},
    {scope: ["X0_0", "X1_0", "X2_0", "X3_0"], "type:" salldiff, "params": {"metric": "var
\rightarrow", "cost": 65}},
    {scope: ["X0_1", "X1_1", "X2_1", "X3_1"], "type:" salldiff, "params": {"metric": "var
\leftrightarrow", "cost": 65}},
    {scope: ["X0_2", "X1_2", "X2_2", "X3_2"], "type:" salldiff, "params": {"metric": "var
   ', "cost": 65}},
    {scope: ["X0_3", "X1_3", "X2_3", "X3_3"], "type:" salldiff, "params": {"metric": "var
\rightarrow", "cost": 65}},
    {scope: ["X0_0"], "costs": [4, 4, 3, 4]},
    {scope: ["X0_1"], "costs": [4, 3, 4, 4]},
    {scope: ["X0_2"], "costs": [2, 1, 3, 2]},
    {scope: ["X0_3"], "costs": [1, 2, 3, 4]},
    {scope: ["X1_0"], "costs": [3, 1, 3, 3]},
    {scope: ["X1_1"], "costs": [4, 1, 1, 1]},
    {scope: ["X1_2"], "costs": [4, 1, 1, 3]},
    {scope: ["X1_3"], "costs": [4, 4, 1, 4]},
{scope: ["X2_0"], "costs": [1, 3, 3, 2]},
    {scope: ["X2_1"], "costs": [2, 1, 3, 1]},
```

	{scope:	["X2_2"],	"costs":	[3,	4,	2,	2]},
	{scope:	["X2_3"],	"costs":	[2,	3,	1,	3]},
	{scope:	["X3_0"],	"costs":	[3,	4,	4,	2]},
	{scope:	["X3_1"],	"costs":	[3,	2,	4,	4]},
	{scope:	["X3_2"],	"costs":	[4,	1,	3,	4]},
	{scope:	["X3_3"],	"costs":	[4,	4,	4,	3]}}
}							

Optimal solution with cost 35 for the latin 4-square example (in red, costs associated to the selected values) :

4, 4, 3, <mark>4</mark>	4, 3, <b>4</b> , 4	2, 1, 3, 2	1, <mark>2</mark> , 3, 4
<b>3</b>	<b>2</b>	0	1
3, <b>1</b> , 3, 3	4, 1, 1, <b>1</b>	4, 1, <b>1</b> , 3	4, 4, 1, 4
1	3	<b>2</b>	0
1, 3, 3, 2	2, <b>1</b> , 3, 1	3, 4, 2, <mark>2</mark>	2, 3, <b>1</b> , 3
0	1	3	<b>2</b>
3, 4, <b>4</b> , 2	3, 2, 4, 4	4, <mark>1</mark> , 3, 4	4, 4, 4, <mark>3</mark>
<b>2</b>	0	1	

#### 2.4.4 Python model

The following code using the pytoulbar2 library solves the weighted latin square problem with the first argument being the dimension N of the chessboard (e.g. "python3 latinsquare.py 6").

```
latinsquare.py
```

```
import sys
from random import seed, randint
seed(123456789)
import pytoulbar2
N = int(sys.argv[1])
top = N**3 +1
Problem = pytoulbar2.CFN(top)
for i in range(N):
    #Create a variable for each square
    Problem.AddVariable('Cell(' + str(i) + ',' + str(j) + ')', range(N))
for i in range(N):
    #Create a constraint all different with variables on the same row
    Problem.AddAllDifferent(['Cell(' + str(i) + ',' + str(j) + ')' for j in range(N)],_
```

```
→encoding = 'salldiffkp')
    #Create a constraint all different with variables on the same column
   Problem.AddAllDifferent(['Cell(' + str(j) + ',' + str(i) + ')'for j in range(N)],__
→encoding = 'salldiffkp')
#Random unary costs
for i in range(N):
   for j in range(N):
        ListConstraintsUnaryC = []
        for 1 in range(N):
            ListConstraintsUnaryC.append(randint(1,N))
        Problem.AddFunction(['Cell(' + str(i) + ', ' + str(j) + ')'],_
→ListConstraintsUnaryC)
#Problem.Dump('WeightLatinSquare.cfn')
Problem.CFN.timer(300)
res = Problem.Solve(showSolutions = 3)
if res and len(res[0]) == N*N:
    # pretty print solution
   for i in range(N):
        print([res[0][i * N + j] for j in range(N)])
    # and its cost
   print("Cost:", int(res[1]))
```

#### 2.4.5 C++ model

The following code using the C++ toulbar2 library API solves the weighted latin square problem.

latinsquare.cpp

```
#include <iostream>
#include <iostream>
#include <vector>
#include "core/tb2wcsp.hpp"
using namespace std;
// an alias for storing the variable costs
// first dim is the grid rows and second is the columns
typedef std::vector<std::vector<cost>>> LatinCostArray;
/*!
    \brief generate random costs for each variable (cell)
*/
void initLatinCosts(size_t N, LatinCostArray& costs) {
    // N*N*N values, costs for each cell
    costs.resize(N);
    for(auto& col: costs) {
        col.resize(N);
    }
}
```

```
for(auto& cell: col) {
             cell.resize(N);
             for(size_t val_ind = 0; val_ind < N; val_ind += 1) {</pre>
                 cell[val_ind] = (rand()%N)+1;
             }
        }
    }
}
/*!
    \brief print the costs for each unary variabl (cell)
*/
void printCosts(LatinCostArray& costs) {
    for(size_t row_ind = 0; row_ind < costs.size(); row_ind ++) {</pre>
        for(size_t col_ind = 0; col_ind < costs[row_ind].size(); col_ind ++) {</pre>
             cout << "cell " << row_ind << "_" << col_ind;</pre>
             cout << " : ";</pre>
             for(auto& cost: costs[row_ind][col_ind]) {
                 cout << cost << ", ";</pre>
             }
            cout << endl;</pre>
        }
    }
}
/*!
    \brief fill in a WCSP object with a latin square problem
 */
void buildWCSP(WeightedCSP& wcsp, LatinCostArray& costs, size_t N, Cost top) {
    // variables
    for(size_t row = 0; row < N; row ++) {</pre>
        for(unsigned int col = 0; col < N; col ++) {</pre>
             wcsp.makeEnumeratedVariable("Cell_" + to_string(row) + "," + to_string(col),_
\rightarrow 0, N-1);
        }
    }
    cout << "number of variables: " << wcsp.numberOfVariables() << endl;</pre>
    /* costs for all different constraints (top on diagonal) */
    vector<Cost> alldiff_costs;
    for(unsigned int i = 0; i < N; i ++) {
        for(unsigned int j = 0; j < N; j ++) {
             if(i == j) {
                 alldiff_costs.push_back(top);
             } else {
                 alldiff_costs.push_back(0);
             }
        }
    }
```

```
/* all different constraints */
    for(unsigned int index = 0; index < N; index ++) {</pre>
        for(unsigned int var_ind1 = 0; var_ind1 < N; var_ind1 ++) {</pre>
            for(unsigned int var_ind2 = var_ind1+1; var_ind2 < N; var_ind2 ++) {</pre>
                 /* row constraints */
                 wcsp.postBinaryConstraint(N*index+var_ind1, N*index+var_ind2, alldiff_
→costs);
                 /* col constraints */
                 wcsp.postBinaryConstraint(index+var_ind1*N, index+var_ind2*N, alldiff_
\rightarrowcosts);
            }
        }
    }
    /* unary costs */
    size_t var_ind = 0;
    for(size_t row = 0; row < N; row ++) {</pre>
        for(size_t col = 0; col < N; col ++) {</pre>
            wcsp.postUnaryConstraint(var_ind, costs[row][col]);
            var_ind += 1;
        }
    }
}
int main() {
    srand(123456789);
    size_t N = 5;
    Cost top = N*N*N + 1;
    // N*N*N values, costs for each cell
    LatinCostArray objective_costs;
    // init the costs for each cell
    initLatinCosts(N, objective_costs);
    cout << "Randomly genereated costs : " << endl;</pre>
    printCosts(objective_costs);
    cout << endl;</pre>
    tb2init();
    ToulBar2::verbose = 0;
    WeightedCSPSolver* solver = WeightedCSPSolver::makeWeightedCSPSolver(top);
    // fill in the WeightedCSP object
    WeightedCSP* wcsp = solver->getWCSP();
```

```
buildWCSP(*wcsp, objective_costs, N, top);
              bool result = solver->solve();
              if(result) {
                              Cost bestCost = solver->getSolutionValue();
                              Cost bestLowerBound = solver->getDDualBound();
                              if(!ToulBar2::limited) {
                                             cout << "Optimal solution found with cost " << bestCost << endl;</pre>
                              } else {
                                             cout << "Best solution found with cost " << bestCost << " and best lower_

where the set of the set of
                              }
                              // retrieve the solution
                              std::vector<Value> solution = solver->getSolution();
                              cout << endl << "Best solution : " << endl;</pre>
                              for(size_t var_ind = 0; var_ind < solution.size(); var_ind ++) {</pre>
                                             cout << solution[var_ind] << " ; ";</pre>
                                             if((var_ind+1) % N == 0) {
                                                             cout << endl;</pre>
                                             }
                              }
              } else {
                              cout << "No solution has been found !" << endl;</pre>
               }
              delete solver;
              return 0:
}
```

## 2.5 Bicriteria weighted latin square problem

#### 2.5.1 Brief description

In this variant of the *Weighted latin square problem*, the objective (sum of the costs of the cells) is decomposed into two criteria: the sum of the cells in the first half of the chessboard and the sum of the cells in the second half. A subset of the pareto solutions can be obtained by solving linear combinations of the two criteria with various weights on the objectives. This can be achieved in ToulBar2 via a MultiCFN object.

#### 2.5.2 CFN model

Similarly to the *Weighted latin square problem*, NxN variables are created with a domain size N. In this model, the permutation of every row and every column is ensured through infinite costs in binary cost functions. Two different CFN are created to represent the two objectives: a first CFN where unary costs are added only for the first half of the chessboard, and a second one with unary costs for the remaining cells.

Toulbar2 allows to either solve for a chosen weighted sum of the two cost function networks as input, or approximate the pareto front by enumerating a complete set of non-redundant weights. As it is shown below, the method allows to compute solutions which costs lie in the convex hull of the pareto front. However, potential solutions belonging to the triangles will be missed with this approach.



#### 2.5.3 Python model

The following code using the pytoulbar2 library solves the bicriteria weighted latin square problem with two different pairs of weights for the two objectives.

bicriteria\_latinsquare.py

```
import sys
from random import seed, randint
seed(123456789)
import pytoulbar2
from matplotlib import pyplot as plt
N = int(sys.argv[1])
top = N^{**3} + 1
# printing a solution as a grid
def print_solution(sol, N):
  grid = [0 for _ in range(N*N)]
  for k,v in sol.items():
    grid[ int(k[5])*N+int(k[7]) ] = int(v[1:])
  output = ''
  for var_ind in range(len(sol)):
    output += str(grid[var_ind]) + ' '
    if var_ind % N == N-1:
      output += ' n'
```

```
print(output, end='')
# creation of the base problem: variables and hard constraints (alldiff must be
→ decomposed into binary constraints)
def create_base_cfn(cfn, N, top):
  # variable creation
 var_indexes = []
  # create N^2 variables, with N values in their domains
  for row in range(N):
   for col in range(N):
      index = cfn.AddVariable('Cell_' + str(row) + '_' + str(col), ['v' + str(val) for_
→val in range(N)])
     var_indexes.append(index)
  # all permutation constraints: pairwise all different
  # forbidden values are enforced by infinite costs
  alldiff_costs = [ top if row == col else 0 for row in range(N) for col in range(N) ]
  for index in range(N):
   for var_ind1 in range(N):
      for var_ind2 in range(var_ind1+1, N):
        # permutations in the rows
        cfn.AddFunction([var_indexes[N*index+var_ind1], var_indexes[N*index+var_ind2]],
→alldiff_costs)
        # permutations in the columns
       cfn.AddFunction([var_indexes[index+var_ind1*N], var_indexes[index+var_ind2*N]],
→alldiff_costs)
split_index = (N*N)//2
# generation of random costs
cell_costs = [[randint(1,N) for _ in range(N)] for _ in range(N*N)]
# multicfn is the main object for combining multiple cost function networks
multicfn = pytoulbar2.MultiCFN()
# first cfn: first half of the grid
cfn = pytoulbar2.CFN(ubinit = top, resolution=6)
cfn.SetName('first half')
create_base_cfn(cfn, N, top)
for variable_index in range(split_index):
  cfn.AddFunction([variable_index], cell_costs[variable_index])
multicfn.PushCFN(cfn)
```

```
# second cfn: second half of the grid
cfn = pytoulbar2.CFN(ubinit = top, resolution=6)
cfn.SetName('second half')
create_base_cfn(cfn, N, top)
for variable_index in range(split_index+1, N*N):
  cfn.AddFunction([variable_index], cell_costs[variable_index])
multicfn.PushCFN(cfn)
# solve with a first pair of weights
weights = (1., 2.)
multicfn.SetWeight(0, weights[0])
multicfn.SetWeight(1, weights[1])
cfn = pytoulbar2.CFN()
cfn.InitFromMultiCFN(multicfn) # the final cfn is initialized from the combined cfn
# cfn.Dump('python_latin_square_bicriteria.cfn')
result = cfn.Solve(timeLimit = 60)
if result:
 print('Solution found with weights', weights, ':')
  sol_costs = multicfn.GetSolutionCosts()
  solution = multicfn.GetSolution()
  print_solution(solution, N)
 print('with costs:', sol_costs, '(weighted sum=', result[1], ')')
print('\n')
# solve a second time with other weights
weights = (2.5, 1.)
multicfn.SetWeight(0, weights[0])
multicfn.SetWeight(1, weights[1])
cfn = pytoulbar2.CFN()
cfn.InitFromMultiCFN(multicfn) # the final cfn is initialized from the combined cfn
# cfn.Dump('python_latin_square_bicriteria.cfn')
result = cfn.Solve(timeLimit = 60)
if result:
 print('Solution found with weights', weights, ':')
  sol_costs = multicfn.GetSolutionCosts()
  solution = multicfn.GetSolution()
  print_solution(solution, N)
  print('with costs:', sol_costs, '(weighted sum=', result[1], ')')
```

```
# approximate the pareto front
(solutions, costs) = multicfn.ApproximateParetoFront(0, 'min', 1, 'min', showSolutions =_
\leftrightarrow 0, timeLimit = 300, timeLimit_per_solution = 60)
fig, ax = plt.subplots()
ax.scatter([c[0]  for c in costs], [c[1]  for c in costs], marker='x')
for index in range(len(costs)-1):
 ax.plot([costs[index][0], costs[index+1][0]], [costs[index][1],costs[index+1][1]], '--
\rightarrow', c='k')
 ax.plot([costs[index][0], costs[index+1][0]], [costs[index][1], costs[index][1]], '--',
\rightarrow c='red')
  ax.plot([costs[index+1][0], costs[index+1][0]], [costs[index][1],costs[index+1][1]], '-
\hookrightarrow-', c='red')
ax.set_xlabel('First objective')
ax.set_ylabel('Second objective')
ax.set_title('Approximation of the Pareto front')
ax.set_aspect('equal')
plt.grid()
plt.show()
```

#### 2.5.4 C++ model

The following code using the C++ toulbar2 library API solves the weighted latin square problem.

```
bicriteria_latinsquare.cpp
```

```
#include <iostream>
#include <vector>
#include "core/tb2wcsp.hpp"
#include "mcriteria/multicfn.hpp"
#include "mcriteria/bicriteria.hpp"
using namespace std;
// an alias for storing the variable costs
// first dim is the grid rows and second is the columns
typedef std::vector<std::vector<Cost>>> LatinCostArray;
// generate random costs for each variable (cell)
// param N grid size
// param costs the matrix costs
void createCostMatrix(size_t N, LatinCostArray& costs) {
   // N*N*N values, costs for each cell
   costs.resize(N);
   for(auto& col: costs) {
       col.resize(N);
       for(auto& cell: col) {
```

```
cell.resize(N);
            for(size_t val_ind = 0; val_ind < N; val_ind += 1) {</pre>
                 cell[val_ind] = (rand()%N)+1;
            }
        }
    }
}
// print the costs for each unary variabl (cell)
// param costs the cost matrix
void printCosts(LatinCostArray& costs) {
    for(size_t row_ind = 0; row_ind < costs.size(); row_ind ++) {</pre>
        for(size_t col_ind = 0; col_ind < costs[row_ind].size(); col_ind ++) {</pre>
            cout << "cell " << row_ind << "_" << col_ind;</pre>
            cout << " : ";</pre>
            for(auto& cost: costs[row_ind][col_ind]) {
                 cout << cost << ", ";</pre>
            }
            cout << endl;</pre>
        }
    }
}
// fill in a WCSP object with a latin square problem
// param wcsp the wcsp object to fill
// param LatinCostArray the cost matrix
// param N grid size
// top the top value, problem upper bound (the objective is always lower than top)
void buildLatinSquare(WeightedCSP& wcsp, LatinCostArray& costs, size_t N, Cost top) {
    // variables
    for(size_t row = 0; row < N; row ++) {</pre>
        for(unsigned int col = 0; col < N; col ++) {</pre>
            wcsp.makeEnumeratedVariable("Cell_" + to_string(row) + "," + to_string(col),_
\rightarrow 0, N-1);
        }
    }
    /* costs for all different constraints (top on diagonal) */
    vector<Cost> alldiff_costs;
    for(unsigned int i = 0; i < N; i ++) {
        for(unsigned int j = 0; j < N; j ++) {
            if(i == j) {
                 alldiff_costs.push_back(top);
            } else {
                 alldiff_costs.push_back(0);
            }
        }
    }
```

```
/* all different constraints */
    for(unsigned int index = 0; index < N; index ++) {</pre>
        for(unsigned int var_ind1 = 0; var_ind1 < N; var_ind1 ++) {</pre>
            for(unsigned int var_ind2 = var_ind1+1; var_ind2 < N; var_ind2 ++) {</pre>
                 /* row constraints */
                 wcsp.postBinaryConstraint(N*index+var_ind1, N*index+var_ind2, alldiff_
→costs);
                 /* col constraints */
                 wcsp.postBinaryConstraint(index+var_ind1*N, index+var_ind2*N, alldiff_
\rightarrowcosts);
            }
        }
    }
    /* unary costs */
    size_t var_ind = 0;
    for(size_t row = 0; row < N; row ++) {</pre>
        for(size_t col = 0; col < N; col ++) {</pre>
            wcsp.postUnaryConstraint(var_ind, costs[row][col]);
            var_ind += 1;
        }
    }
}
// print a solution as a grid
// param N the size of the grid
// param solution the multicfn solution (dict)
// param point the objective costs (objective space point)
void printSolution(size_t N, MultiCFN::Solution& solution, Bicriteria::Point& point) {
    for(size_t row = 0; row < N; row ++) {</pre>
        for(size_t col = 0; col < N; col ++) {</pre>
            string var_name = "Cell_" + to_string(row) + "," + to_string(col);
            cout << solution[var_name].substr(1) << " ";</pre>
        }
        cout << endl;</pre>
    }
    cout << "obj_1 = " << point.first << "; obj2 = " << point.second << endl;</pre>
}
// main function
int main() {
    srand(123456789);
    size_t N = 4;
    Cost top = N*N*N + 1;
    // two cost matrice
    LatinCostArray costs_obj1, costs_obj2;
```

```
// init the objective with random costs
   createCostMatrix(N, costs_obj1);
   createCostMatrix(N, costs_obj2);
   // cout << "Randomly genereated costs : " << endl;</pre>
   // printCosts(costs_obj1);
   // cout << endl << endl;</pre>
   // printCosts(costs_obj2);
   tb2init();
   initCosts();
    // create the two wcsp objects
   WeightedCSP* wcsp1 = WeightedCSP::makeWeightedCSP(top);
   WeightedCSP* wcsp2 = WeightedCSP::makeWeightedCSP(top);
   // initialize the objects as a latin square problem objectives with two different.
→objectves
   buildLatinSquare(*wcsp1, costs_obj1, N, top);
   buildLatinSquare(*wcsp2, costs_obj2, N, top);
   // creation of the multicfn
   MultiCFN mcfn:
   mcfn.push_back(dynamic_cast<WCSP*>(wcsp1));
   mcfn.push_back(dynamic_cast<WCSP*>(wcsp2));
    // computation iof the supported points of the biobjective problem
   Bicriteria::computeSupportedPoints(&mcfn, std::make_pair(Bicriteria::OptimDir::Optim_

→Min, Bicriteria::OptimDir::Optim_Min));

   // access to the computed solutions and their objective values
    std::vector<MultiCFN::Solution> solutions = Bicriteria::getSolutions();
   std::vector<Bicriteria::Point> points = Bicriteria::getPoints();
   // print all solutions computed
   cout << "Resulting solutions: " << endl;</pre>
   for(size_t sol_index = 0; sol_index < solutions.size(); sol_index ++) {</pre>
       printSolution(N, solutions[sol_index], points[sol_index]);
        cout << endl;</pre>
   }
   // delete the wcsp objects
   delete wcsp1;
   delete wcsp2;
   return 0;
}
```

The above code can be compiled with the following command:

```
g++ -O3 -std=c++17 -Wall -DBOOST -DLONGLONG_COST -DLONGDOUBLE_PROB -I $YOUR_TB2_INCLUDE_
(continues on next page)
```

→PATH main.cpp -c -o main.o

Where **\$YOUR\_TB2\_INCLUDE\_PATH** is the path to the ToulBar2 src directory. And the compiled program is obtained via :

```
g++ -O3 -std=c++17 -Wall -DBOOST -DLONGLONG_COST -DLONGDOUBLE_PROB main.o -o main -L

→$YOUR_LIBTB2_PATH -ltb2 -lgmp -lboost_graph -lboost_iostreams -lz -llzma
```

Where **\$YOUR\_LIBTB2\_PATH** is the path to the ToulBar2 compiled library. When running the program, do not forget to set the **\$(LD\_LIBRARY\_PATH)** environment variable in Linux.

## 2.6 Radio link frequency assignment problem

#### 2.6.1 Brief description

The problem consists in assigning frequencies to radio communication links in such a way that no interferences occur. Domains are set of integers (non-necessarily consecutive).

Two types of constraints occur:

- (I) the absolute difference between two frequencies should be greater than a given number  $d_i (|x y| > d_i)$
- (II) the absolute difference between two frequencies should exactly be equal to a given number  $d_i (|x y| = d_i)$ .

Different deviations  $d_i$ , i in 0..4, may exist for the same pair of links.  $d_0$  corresponds to hard constraints while higher deviations are soft constraints that can be violated with an associated cost  $a_i$ . Moreover, pre-assigned frequencies may be known for some links which are either hard or soft preferences (mobility cost  $b_i$ , i in 0..4). The goal is to minimize the weighted sum of violated constraints.

#### So the goal is to minimize the sum:

a\_1\*nc1+...+a\_4\*nc4+b\_1\*nv1+...+b\_4\*nv4

where nci is the number of violated constraints with cost a\_i and nvi is the number of modified variables with mobility cost b\_i.

Cabon, B., de Givry, S., Lobjois, L., Schiex, T., Warners, J.P. Constraints (1999) 4: 79.

#### 2.6.2 CFN model

We create N variables for every radio link with a given integer domain. Hard and soft binary cost functions express interference constraints with possible deviations with cost equal to  $a_i$ . Unary cost functions are used to model mobility costs with cost equal to  $b_i$ . The initial upper bound is defined as 1 plus the total cost where all the soft constraints are maximally violated (costs  $a_4/b_4$ ).

#### 2.6.3 Data

Original data files can be downloaded from the cost function library FullRLFAP. Their format is described here. You can try a small example CELAR6-SUB1 (var.txt, dom.txt, ctr.txt, cst.txt) with optimum value equal to 2669.

#### 2.6.4 Python model

The following code solves the corresponding cost function network using the pytoulbar2 library and needs 4 arguments: the variable file, the domain file, the constraints file and the cost file (e.g. "python3 rlfap.py var.txt dom.txt ctr.txt cst.txt").

rlfap.py

```
import svs
import pytoulbar2
class Data:
        def __init__(self, var, dom, ctr, cst):
                self.var = list()
                self.dom = {}
                self.ctr = list()
                self.cost = {}
                self.nba = {}
                self.nbb = {}
                self.top = 1
                self.Domain = {}
                stream = open(var)
                for line in stream:
                         if len(line.split())>=4:
                                 (varnum, vardom, value, mobility) = line.split()[:4]
                                 self.Domain[int(varnum)] = int(vardom)
                                 self.var.append((int(varnum), int(vardom), int(value),__

→int(mobility)))

                                 self.nbb["b" + str(mobility)] = self.nbb.get("b" +_
\rightarrow str(mobility), \emptyset) + 1
                         else:
                                 (varnum, vardom) = line.split()[:2]
                                 self.Domain[int(varnum)] = int(vardom)
                                 self.var.append((int(varnum), int(vardom)))
                stream = open(dom)
                for line in stream:
                         domain = line.split()[:]
                         self.dom[int(domain[0])] = [int(f) for f in domain[2:]]
                stream = open(ctr)
                for line in stream:
                         (var1, var2, dummy, operand, deviation, weight) = line.
\rightarrow split()[:6]
                         self.ctr.append((int(var1), int(var2), operand, int(deviation),__
→int(weight)))
                         self.nba["a" + str(weight)] = self.nba.get("a" + str(weight), 0)_
<u>→</u>+ 1
                stream = open(cst)
                for line in stream:
                         if len(line.split()) == 3:
                                 (aorbi, eq, cost) = line.split()[:3]
                                 if (eq == "="):
                                          self.cost[aorbi] = int(cost)
                                          self.top += int(cost) * self.nba.get(aorbi, self.

→nbb.get(aorbi, 0))

#collect data
data = Data(sys.argv[1], sys.argv[2], sys.argv[3], sys.argv[4])
```

```
top = data.top
Problem = pytoulbar2.CFN(top)
#create a variable for each link
for e in data.var:
        domain = []
        for f in data.dom[e[1]]:
                domain.append('f' + str(f))
        Problem.AddVariable('link' + str(e[0]), domain)
#binary hard and soft constraints
for (var1, var2, operand, deviation, weight) in data.ctr:
        ListConstraints = []
        for a in data.dom[data.Domain[var1]]:
                for b in data.dom[data.Domain[var2]]:
                        if ((operand==">" and abs(a - b) > deviation) or (operand=="="_
→and abs(a - b) == deviation)):
                                ListConstraints.append(0)
                        else:
                                ListConstraints.append(data.cost.get('a' + str(weight),
→top))
        Problem.AddFunction(['link' + str(var1), 'link' + str(var2)], ListConstraints)
#unary hard and soft constraints
for e in data.var:
        if len(e) >= 3:
                ListConstraints = [1]
                for a in data.dom[e[1]]:
                        if a == e[2]:
                                ListConstraints.append()
                        else:
                                ListConstraints.append(data.cost.get('b' + str(e[3]),
→top))
                Problem.AddFunction(['link' + str(e[0])], ListConstraints)
#Problem.Dump('rlfap.cfn')
Problem.CFN.timer(300)
res = Problem.Solve(showSolutions=3)
if res:
       print("Best solution found with cost:",int(res[1]),"in", Problem.GetNbNodes(),
→ "search nodes.")
else:
        print('Sorry, no solution found!')
```

## 2.7 Frequency assignment problem with polarization

#### 2.7.1 Brief description

The previously-described *Radio link frequency assignment problem* has been extended to take into account polarization constraints and user-defined relaxation of electromagnetic compatibility constraints. The problem is to assign a pair

(frequency, polarization) to every radio communication link (also called a path). Frequencies are integer values taken in finite domains. Polarizations are in  $\{-1,1\}$ . Constraints are :

- (I) two paths must use equal or different frequencies  $(f_i=f_j \text{ or } f_i <> f_j)$ ,
- (II) the absolute difference between two frequencies should exactly be equal or different to a given number e  $(|f_i f_j| = e \text{ or } |f_i f_j| <> e)$ ,
- (III) two paths must use equal or different polarizations  $(p_i=p_j \text{ or } p_i <> p_j)$ ,
- (IV) the absolute difference between two frequencies should be greater at a relaxation level 1 (0 to 10) than a given number g\_l (resp. d\_l) if polarization are equal (resp. different) ( $|f_i f_j| \ge g_l$  if  $p_i = p_j$  else  $|f_i f_j| \ge d_l$ , with  $g_l l \ge g_l$ ,  $d_l l \ge d_l$ , and usually  $g_l \ge d_l$ .

Constraints (I) to (III) are mandatory constraints, while constraints (IV) can be relaxed. The goal is to find a feasible assignment with the smallest relaxation level 1 and which minimizes the (weighted) number of violations of (IV) at lower levels. See ROADEF\_Challenge\_2001.

The cost of a given solution will be calculated by the following formula: 10\*k\*nbsoft\*\*2 + 10\*nbsoft\*V(k-1) + V(k-2) + V(k-3) + ... + V0

where nbsoft is the number of soft constraints in the problem and k the feasible relaxation level and V(i) the number of violated contraints of level i.



#### 2.7.2 CFN model

We create a single variable to represent a pair (frequency,polarization) for every radio link, but be aware, toulbar2 can only read str or int values, so in order to give a tuple to toulbar2 we need to first transform them into string. We use hard binary constraints to modelize (I) to (III) type constraints.

We assume the relaxation level k is given as input. In order to modelize (IV) type constraints we first take in argument the level of relaxation i, and we create 11 constraints, one for each relaxation level from 0 to 10. The first k-2 constraints are soft and with a violation cost of 1. The soft constraint at level k-1 has a violation cost 10\*nbsoft (the number of soft constraints) in order to maximize first the number of satisfied constraints at level k-1 and then the other soft constraints. The constraints at levels k to 10 are hard constraints.

The initial upper bound of the problem will be  $10^{(k+1)*nbsoft**2+1}$ .

#### 2.7.3 Data

Original data files can be download from ROADEF or fapp. Their format is described here. You can try a small example exemple1.in (resp. exemple2.in) with optimum 523 at relaxation level 3 with 1 violation at level 2 and 3 below (resp. 13871 at level 7 with 1 violation at level 6 and 11 below). See ROADEF Challenge 2001 results.

#### 2.7.4 Python model

The following code solves the corresponding cost function network using the pytoulbar2 library and needs 4 arguments: the data file and the relaxation level (e.g. "python3 fapp.py exemple1.in 3"). You can also compile fappeval.c using "gcc -o fappeval fappeval.c" and download sol2fapp.awk in order to evaluate the solutions (e.g., "python3 fapp.py exemple1.in 3 | awk -f ./sol2fapp.awk - exemple1").

fapp.py

```
import sys
import pytoulbar2
class Data:
        def __init__(self, filename, k):
                self.var = {}
                self.dom = {}
                self.ctr = list()
                self.softeg = list()
                self.softne = list()
                self.nbsoft = 0
                stream = open(filename)
                for line in stream:
                        if len(line.split())==3 and line.split()[0]=="DM":
                                 (DM, dom, freq) = line.split()[:3]
                                 if self.dom.get(int(dom)) is None:
                                         self.dom[int(dom)] = [int(freq)]
                                 else:
                                         self.dom[int(dom)].append(int(freq))
                        if len(line.split()) == 4 and line.split()[0]=="TR":
                                 (TR, route, dom, polarisation) = line.split()[:4]
                                 if int(polarisation) == 0:
                                         self.var[int(route)] = [(f,-1) for f in self.

    dom[int(dom)]] + [(f,1) for f in self.dom[int(dom)]]

                                 if int(polarisation) == -1:
                                         self.var[int(route)] = [(f,-1) for f in self.
→dom[int(dom)]]
                                 if int(polarisation) == 1:
                                         self.var[int(route)] = [(f,1) for f in self.
→dom[int(dom)]]
                        if len(line.split())==6 and line.split()[0]=="CI":
                                 (CI, route1, route2, vartype, operator, deviation) =
\rightarrow line.split()[:6]
                                 self.ctr.append((int(route1), int(route2), vartype,__
→operator, int(deviation)))
                        if len(line.split())==14 and line.split()[0]=="CE":
                                 (CE, route1, route2, s0, s1, s2, s3, s4, s5, s6, s7, s8,
\rightarrow s9, s10) = line.split()[:14]
                                 self.softeq.append((int(route1), int(route2), [int(s)]
→ for s in [s0, s1, s2, s3, s4, s5, s6, s7, s8, s9, s10]]))
                                 self.nbsoft += 1
```

```
if len(line.split())==14 and line.split()[0]=="CD":
                                 (CD, route1, route2, s0, s1, s2, s3, s4, s5, s6, s7, s8,
\rightarrow s9, s10) = line.split()[:14]
                                 self.softne.append((int(route1), int(route2), [int(s)]
→for s in [s0, s1, s2, s3, s4, s5, s6, s7, s8, s9, s10]]))
                self.top = 10*(k+1)*self.nbsoft**2 + 1
if len(sys.argv) < 2:</pre>
        exit('Command line argument is composed of the problem data filename and the
→relaxation level')
k = int(sys.argv[2])
#collect data
data = Data(sys.argv[1], k)
Problem = pytoulbar2.CFN(data.top)
#create a variable for each link
for e in list(data.var.keys()):
        domain = []
        for i in data.var[e]:
                domain.append(str(i))
        Problem.AddVariable("X" + str(e), domain)
#hard binary constraints
for (route1, route2, vartype, operand, deviation) in data.ctr:
        Constraint = []
        for (f1,p1) in data.var[route1]:
                 for (f2,p2) in data.var[route2]:
                         if vartype == 'F':
                                  if operand == 'E':
                                          if abs(f2 - f1) == deviation:
                                                  Constraint.append(0)
                                          else:
                                                  Constraint.append(data.top)
                                  else:
                                          if abs(f2 - f1) != deviation:
                                                  Constraint.append(♥)
                                          else:
                                                  Constraint.append(data.top)
                         else:
                                  if operand == 'E':
                                          if p2 == p1:
                                                  Constraint.append(♥)
                                          else:
                                                  Constraint.append(data.top)
                                  else:
                                          if p2 != p1:
                                                  Constraint.append(0)
                                          else:
```

```
Constraint.append(data.top)
        Problem.AddFunction(["X" + str(route1), "X" + str(route2)], Constraint)
#soft binary constraints for equal polarization
for (route1, route2, deviations) in data.softeq:
        for i in range(11):
                ListConstraints = []
                for (f1,p1) in data.var[route1]:
                        for (f2,p2) in data.var[route2]:
                                  if p1!=p2 or abs(f1 - f2) >= deviations[i]:
                                          ListConstraints.append(0)
                                  elif i >= k:
                                          ListConstraints.append(data.top)
                                  elif i == k-1:
                                          ListConstraints.append(10*data.nbsoft)
                                  else:
                                          ListConstraints.append(1)
                Problem.AddFunction(["X" + str(route1), "X" + str(route2)],_
→ListConstraints)
#soft binary constraints for not equal polarization
for (route1, route2, deviations) in data.softne:
        for i in range(11):
                ListConstraints = []
                for (f1,p1) in data.var[route1]:
                         for (f2,p2) in data.var[route2]:
                                  if p1==p2 or abs(f1 - f2) >= deviations[i]:
                                          ListConstraints.append(0)
                                  elif i >= k:
                                          ListConstraints.append(data.top)
                                  elif i == k-1:
                                          ListConstraints.append(10*data.nbsoft)
                                  else:
                                          ListConstraints.append(1)
                Problem.AddFunction(["X" + str(route1), "X" + str(route2)],_
→ListConstraints)
#zero-arity cost function representing a constant cost corresponding to the relaxation...
\rightarrow at level k
Problem.AddFunction([], 10*k*data.nbsoft**2)
#Problem.Dump('Fapp.cfn')
Problem.CFN.timer(900)
Problem.Solve(showSolutions=3)
```

## 2.8 Mendelian error detection problem

#### 2.8.1 Brief description

The problem is to detect marker genotyping incompatibilities (Mendelian errors) in complex pedigrees. The input is a pedigree data with partial observed genotyping data at a single locus, we assume the pedigree to be exact, but not the genotyping data. The problem is to assign genotypes (unordered pairs of alleles) to all individuals such that they are compatible with the Mendelian law of heredity (one allele is the same as their father's and one as their mother's). The goal is to maximize the number of matching alleles between the genotyping data and the solution. Each difference from the genotyping data has a cost of 1.

Sanchez, M., de Givry, S. and Schiex, T. Constraints (2008) 13:130.

#### 2.8.2 CFN model

We create N variables, one for each individual genotype with domain being all possible unordered pairs of existing alleles. Hard ternary cost functions express mendelian law of heredity (one allele is the same as their father's and one as their mother's, with mother and father defined in the pedigree data). For each genotyping data, we create one unary soft constraint with violation cost equal to 1 to represent the matching between the genotyping data and the solution.

#### 2.8.3 Data

Original data files can be download from the cost function library pedigree. Their format is described here. You can try a small example simple.pre (simple.pre) with optimum value equal to 1.

#### 2.8.4 Python model

The following code solves the corresponding cost function network using the pytoulbar2 library (e.g. "python3 mendel.py simple.pre").

mendel.py

```
import sys
import pytoulbar2
class Data:
        def __init__(self, ped):
                self.id = list()
                self.father = {}
                self.mother = {}
                self.allelesId = {}
                self.ListAlle = list()
                self.obs = 0
                stream = open(ped)
                for line in stream:
                         (locus, id, father, mother, sex, allele1, allele2) = line.
→split()[:]
                        self.id.append(int(id))
                        self.father[int(id)] = int(father)
                        self.mother[int(id)] = int(mother)
                        self.allelesId[int(id)] = (int(allele1), int(allele2)) if_
int(allele1) < int(allele2) else (int(allele2), int(allele1))</pre>
                         if not(int(allele1) in self.ListAlle) and int(allele1) != 0:
                                 self.ListAlle.append(int(allele1))
```

```
(continued from previous page)
                         if int(allele2) != 0 and not(int(allele2) in self.ListAlle):
                                 self.ListAlle.append(int(allele2))
                         if int(allele1) != 0 or int(allele2) != 0:
                                 self.obs += 1
#collect data
data = Data(sys.argv[1])
top = int(data.obs+1)
Problem = pytoulbar2.CFN(top)
#create a variable for each individual
for i in data.id:
        domains = []
        for a1 in data.ListAlle:
                for a2 in data.ListAlle:
                         if a1 <= a2:
                                 domains.append('a'+str(a1)+'a'+str(a2))
        Problem.AddVariable('g' + str(i) , domains)
#create the constraints that represent the mendel's laws
ListConstraintsMendelLaw = []
for p1 in data.ListAlle:
        for p2 in data.ListAlle:
                if p1 <= p2:
                                     # father alleles
                         for m1 in data.ListAlle:
                                 for m2 in data.ListAlle:
                                         if m1 <= m2:
                                                               # mother alleles
                                                  for a1 in data.ListAlle:
                                                           for a2 in data.ListAlle:
                                                                   if a1 <= a2:
                                                                                        #.
→ child alleles
                                                                           if (a1 in (p1,
\rightarrowp2) and a2 in (m1,m2)) or (a2 in (p1,p2) and a1 in (m1,m2)) :
                                                                                    ListConstraintsMendelLa
\rightarrow append(\emptyset)
                                                                           else :
                                                                           ш.
          ListConstraintsMendelLaw.append(top)
for i in data.id:
        #ternary constraints representing mendel's laws
        if data.father.get(i, 0) != 0 and data.mother.get(i, 0) != 0:
                Problem.AddFunction(['g' + str(data.father[i]),'g' + str( data.
→mother[i]), 'g' + str(i)], ListConstraintsMendelLaw)
        #unary constraints linked to the observations
        if data.allelesId[i][0] != 0 and data.allelesId[i][1] != 0:
                ListConstraintsObservation = []
                for a1 in data.ListAlle:
                         for a2 in data.ListAlle:
                                 if a1 <= a2:
                                                                              (continues on next page)
```

```
if (a1,a2) == data.allelesId[i]:
ListConstraintsObservation.append(0)
else :
ListConstraintsObservation.append(1)
Problem.AddFunction(['g' + str(i)], ListConstraintsObservation)
#Problem.CFN.timer(300)
res = Problem.Solve(showSolutions=3)
if res:
    print('There are',int(res[1]),'difference(s) between the solution and the_
->observation.')
else:
    print('No solution found')
```

## 2.9 Block modeling problem

#### 2.9.1 Brief description

This is a clustering problem, occurring in social network analysis.

The problem is to divide a given directed graph G into k clusters such that the interactions between clusters can be summarized by a k k 0/1 matrix M: if M[i,j]=1 then all the nodes in cluster i should be connected to all the nodes in cluster j in G, else if M[i,j]=0 then there should be no edge in G between the nodes from the two clusters.

For example, the following graph G is composed of 4 nodes:



and corresponds to the following matrix:

<u>/</u> 0	0	0	0)
1	0	1	1
0	0	0	0
0	0	0	0/

It can be perfectly clusterized into the following graph by clustering together the nodes 0, 2 and 3 in cluster 1 and the node 1 in cluster 0:



and this graph corresponds to the following M matrix:

 $\begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}$ 

On the contrary, if we decide to cluster the next graph G' in the same way as above, the edge (2, 3) will be 'lost' in the process and the cost of the solution will be 1.



The goal is to find a k-clustering of a given graph and the associated matrix M that minimizes the number of erroneous edges.

A Mattenet, I Davidson, S Nijssen, P Schaus. Generic Constraint-Based Block Modeling Using Constraint Programming. CP 2019, pp656-673, Stamford, CT, USA.

#### 2.9.2 CFN model

We create N variables, one for every node of the graph, with domain size k representing the clustering. We add k\*k Boolean variables for representing M.

For all triplets of two nodes u, v, and one matrix cell M[i,j], we have a ternary cost function that returns a cost of 1 if node u is assigned to cluster i, v to j, and M[i,j]=1 but (u,v) is not in G, or M[i,j]=0 and (u,v) is in G. In order to break symmetries, we constrain the first k-1 node variables to be assigned to a cluster number less than or equal to their index

#### 2.9.3 Data

You can try a small example simple.mat with optimum value equal to 0 for 3 clusters.

Perfect solution for the small example with k=3 (Mattenet et al, CP 2019)



More examples with 3 clusters (Stochastic Block Models [Funke and Becker, Plos One 2019])



See other examples, such as PoliticalActor and more, here : 100.mat | 150.mat | 200.mat | 30.mat | 50.mat | hartford\_drug.mat | kansas.mat | politicalactor.mat | sharpstone.mat | transatlantic.mat.

#### 2.9.4 Python model

The following code using pytoulbar2 library solves the corresponding cost function network (e.g. "python3 block-model.py simple.mat 3").

blockmodel.py

```
import sys
import pytoulbar2
#read adjency matrix of graph G
Lines = open(sys.argv[1], 'r').readlines()
GMatrix = [[int(e) for e in l.split(' ')] for l in Lines]
N = len(Lines)
Top = N*N + 1
```

```
K = int(sys.argv[2])
#give names to node variables
Var = [(chr(65 + i) if N < 28 else "x" + str(i)) for i in range(N)] # Political actor or
\rightarrow anv instance
    Var = ["ron", "tom", "frank", "boyd", "tim", "john", "jeff", "jay", "sandy", "jerry", "darrin
→","ben","arnie"] # Transatlantic
  Var = ["justin", "harry", "whit", "brian", "paul", "ian", "mike", "jim", "dan", "ray", "cliff
#
→ ", "mason", "roy"] # Sharpstone
    Var = ["Sherrif", "CivilDef", "Coroner", "Attorney", "HighwayP", "ParksRes", "GameFish",
#
→ "KansasDOT", "ArmyCorps", "ArmyReserve", "CrableAmb", "FrankCoAmb", "LeeRescue", "Shawney",
→ "BurlPolice", "LyndPolice", "RedCross", "TopekaFD", "CarbFD", "TopekaRBW"] # Kansas
Problem = pytoulbar2.CFN(Top)
#create a Boolean variable for each coefficient of the M GMatrix
for u in range(K):
    for v in range(K):
        Problem.AddVariable("M_" + str(u) + "_" + str(v), range(2))
#create a domain variable for each node in graph G
for i in range(N):
    Problem.AddVariable(Var[i], range(K))
#general case for each edge in G
for u in range(K):
    for v in range(K):
        for i in range(N):
            for j in range(N):
                if i != j:
                    ListCost = []
                    for m in range(2):
                         for k in range(K):
                             for 1 in range(K):
                                 if (u == k and v == 1 and GMatrix[i][j] != m):
                                     ListCost.append(1)
                                 else:
                                     ListCost.append(0)
                    Problem.AddFunction(["M_" + str(u) + "_" + str(v), Var[i], Var[j]],
→ListCost)
# self-loops must be treated separately as they involves only two variables
for u in range(K):
    for i in range(N):
        ListCost = []
        for m in range(2):
            for k in range(K):
                if (u == k and GMatrix[i][i] != m):
                    ListCost.append(1)
                else:
```

```
ListCost.append(0)
        Problem.AddFunction(["M_" + str(u) + "_" + str(u), Var[i]], ListCost)
# breaking partial symmetries by fixing first (K-1) domain variables to be assigned to a.
\rightarrow cluster number less than or equal to their index
for 1 in range(K-1):
    Constraint = []
    for k in range(K):
        if k > 1:
            Constraint.append(Top)
        else:
            Constraint.append(0)
    Problem.AddFunction([Var[1]], Constraint)
Problem.Dump(sys.argv[1].replace('.mat','.cfn'))
Problem.CFN.timer(300)
res = Problem.Solve(showSolutions = 3)
if res:
    print("M matrix:")
    for u in range(K):
        Line = []
        for v in range(K):
            Line.append(res[0][u*K+v])
        print(Line)
    for k in range(K):
        for i in range(N):
            if res[0][K**2+i] == k:
                print("Node",Var[i],"with index",str(i),"is in cluster",
\rightarrow str(res[0][K**2+i]))
```

## 2.10 Airplane landing problem

#### 2.10.1 Brief description

We consider a single plane's landing runway. Given a set of planes with given target landing time, the objective is to minimize the total weighted deviation from the target landing time for each plane.

There are costs associated with landing either earlier or later than the target landing time for each plane.

Each plane has to land within its predetermined time window. For each pair of planes, there is an additional constraint to enforce that the separation time between those planes is larger than a given number.

J.E. Beasley, M. Krishnamoorthy, Y.M. Sharaiha and D. Abramson. Scheduling aircraft landings - the static case. Transportation Science, vol.34, 2000.

#### 2.10.2 CFN model

We create N variables, one for each plane, with domain value equal to all their possible landing time.

Binary hard cost functions express separation times between pairs of planes. Unary soft cost functions represent the weighted deviation for each plane.

#### 2.10.3 Data

Original data files can be download from the cost function library airland. Their format is described here. You can try a small example airland1.txt with optimum value equal to 700.

#### 2.10.4 Python model solver

The following code uses the pytoulbar2 module to generate the cost function network and solve it (e.g. "python3 airland.py airland1.txt").

airland.py

```
import sys
import pytoulbar2
f = open(sys.argv[1], 'r').readlines()
tokens = []
for 1 in f:
    tokens += l.split()
pos = 
def token():
    global pos, tokens
    if (pos == len(tokens)):
        return None
    s = tokens[pos]
    pos += 1
    return int(float(s))
N = token()
token() # skip freeze time
LT = []
PC = []
ST = []
for i in range(N):
    token() # skip appearance time
# Times per plane: {earliest landing time, target landing time, latest landing time}
    LT.append([token(), token()], token()])
# Penalty cost per unit of time per plane:
# [for landing before target, after target]
    PC.append([token(), token()])
# Separation time required after i lands before j can land
    ST.append([token() for j in range(N)])
top = 99999
Problem = pytoulbar2.CFN(top)
for i in range(N):
    Problem.AddVariable('x' + str(i), range(LT[i][0],LT[i][2]+1))
                                                                             (continues on next page)
```

```
for i in range(N):
   ListCost = []
   for a in range(LT[i][0], LT[i][2]+1):
        if a < LT[i][1]:
            ListCost.append(PC[i][0]*(LT[i][1] - a))
        else:
            ListCost.append(PC[i][1]*(a - LT[i][1]))
   Problem AddFunction([i], ListCost)
for i in range(N):
   for j in range(i+1,N):
        Constraint = []
        for a in range(LT[i][0], LT[i][2]+1):
            for b in range(LT[j][0], LT[j][2]+1):
                if a+ST[i][j]>b and b+ST[j][i]>a:
                    Constraint.append(top)
                else:
                    Constraint.append(0)
        Problem.AddFunction([i, j],Constraint)
#Problem.Dump('airplane.cfn')
Problem NoPreprocessing()
Problem.Solve(showSolutions = 3)
```

## 2.11 Warehouse location problem

#### 2.11.1 Brief description

A company considers opening warehouses at some candidate locations with each of them having a maintenance cost if it is open.

The company controls a set of given stores and each of them needs to take supplies to one of the warehouses, but depending on the warehouse chosen, there will be an additional supply cost.

The objective is to choose which warehouse to open and to divide the stores among the open warehouses in order to minimize the total cost of supply and maintenance costs.

#### 2.11.2 CFN model

We create Boolean variables for the warehouses (i.e., open or not) and integer variables for the stores, with domain size the number of warehouses to represent to which warehouse the store will take supplies.

Hard binary constraints represent that a store cannot take supplies from a closed warehouse. Soft unary constraints represent the maintenance cost of the warehouses. Soft unary constraints represent the store's cost regarding which warehouse to take supplies from.

#### 2.11.3 Data

Original data files can be download from the cost function library warehouses. Their format is described here.

#### 2.11.4 Python model solver

The following code uses the pytoulbar2 module to generate the cost function network and solve it (e.g. "python3 warehouse.py cap44.txt 1" found an optimum value equal to 10349757). Other instances are available here in cfn format.

warehouse.py

```
import sys
import pytoulbar2
uncapacitated = True # if True then do not enforce capacity constraints on warehouses
f = open(sys.argv[1], 'r').readlines()
precision = int(sys.argv[2]) # in [0,9], used to convert cost values from float to.
→integer (by 10**precision)
tokens = []
for 1 in f:
   tokens += l.split()
pos = 0
def token():
   global pos, tokens
   if pos == len(tokens):
       return None
   s = tokens[pos]
   pos += 1
   return s
N = int(token()) # number of warehouses
M = int(token()) # number of stores
top = 1 # sum of all costs plus one
CostW = [] # maintenance cost of warehouses
Capacity = [] # capacity limit of warehouses
for i in range(N):
   Capacity.append(int(token()))
   CostW.append(int(float(token()) * 10.**precision))
top += sum(CostW)
Demand = [] # demand for each store
CostS = [[] for i in range(M)] # supply cost matrix
for j in range(M):
   Demand.append(int(token()))
    for i in range(N):
        CostS[j].append(int(float(token()) * 10.**precision))
```

```
top += sum(CostS[j])
# create a new empty cost function network
Problem = pytoulbar2.CFN(top)
# add warehouse variables
for i in range(N):
    Problem.AddVariable('w' + str(i), range(2))
# add store variables
for j in range(M):
    Problem.AddVariable('s' + str(j), range(N))
# add maintenance costs
for i in range(N):
    Problem.AddFunction([i], [0, CostW[i]])
# add supply costs for each store
for j in range(M):
    Problem AddFunction([N+j], CostS[j])
# add channeling constraints between warehouses and stores
for i in range(N):
    for j in range(M):
        Problem.AddFunction([i, N+j], [(top if (a == 0 and b == i) else 0) for a in_
\rightarrow range(2) for b in range(N)])
# optional: add capacity constraint on each warehouse
if not(uncapacitated):
    for i in range(N):
        Problem.AddGeneralizedLinearConstraint([(N+j, i, min(max(Capacity), Demand[j]))_

→ for j in range(M)], '<=', Capacity[i])
</pre>
#Problem.Dump('warehouse.cfn')
Problem.CFN.timer(300)
res = Problem.Solve(showSolutions=3)
if res:
        print("Best solution found with cost:",int(res[1]),"in", Problem.GetNbNodes(),
\rightarrow "search nodes.")
else:
        print('Sorry, no solution found!')
```

## 2.12 Square packing problem

#### 2.12.1 Brief description

We have N squares of respective size  $1 \times 1$ ,  $2 \times 2$ ,..., NxN. We have to fit them without overlaps into a square of size SxS.

Results up to N=56 are given here.

An optimal solution for 15 squares packed into a 36x36 square (Fig. taken from Takehide Soh)

7		14		15
13	;	12		11
9	4	8	5	10

#### 2.12.2 CFN model

We create an integer variable of domain size (S-i)x(S-i) for each square. The variable represents the position of the top left corner of the square.

The value of a given variable modulo (S-i) gives the x-coordinate, whereas its value divided by (S-i) gives the y-coordinate.

We have hard binary constraints to forbid any overlapping pair of squares.

We make the problem a pure satisfaction problem by fixing the initial upper bound to 1.

#### 2.12.3 Python model

The following code uses the pytoulbar2 library to generate the cost function network and solve it (e.g. "python3 square.py 3 5"). square.py

```
import sys
from random import randint, seed
seed(123456789)
import pytoulbar2
try:
       N = int(sys.argv[1])
        S = int(sys.argv[2])
        assert N <= S
except:
       print('Two integers need to be given as arguments: N and S')
        exit()
#pure constraint satisfaction problem
Problem = pytoulbar2.CFN(1)
#create a variable for each square
for i in range(N):
       Problem.AddVariable('sq' + str(i+1), ['(' + str(1) + ',' + str(j) + ')' for 1 in.

→range(S-i) for j in range(S-i)])
```

```
#binary hard constraints for overlapping squares
for i in range(N):
         for j in range(i+1,N):
                  ListConstraintsOverlaps = []
                  for a in [S*k+1 for k in range(S-i) for 1 in range(S-i)]:
                           for b in [S*m+n for m in range(S-j) for n in range(S-j)]:
                                     #calculating the coordinates of the squares
                                    X_i = a\%S
                                    X_j = b\%S
                                    Y_i = a//S
                                    Y_j = b//S
                                    #calculating if squares are overlapping
                                    if X_i >= X_j :
                                              if X_i - X_j < j+1:
                                                       if Y_i >= Y_j:
                                                                if Y_i - Y_j < j+1:
                                                                         ListConstraintsOverlaps.
\rightarrow append(1)
                                                                else:
                                                                         ListConstraintsOverlaps.
\rightarrow append(\emptyset)
                                                       else:
                                                                if Y_j - Y_i < i+1:
                                                                         ListConstraintsOverlaps.
\rightarrow append(1)
                                                                else:
                                                                         ListConstraintsOverlaps.
\rightarrow append(\emptyset)
                                              else:
                                                       ListConstraintsOverlaps.append(0)
                                    else :
                                              if X_j - X_i < i+1:
                                                       if Y_i >= Y_j:
                                                                if Y_i - Y_j < j+1:
                                                                         ListConstraintsOverlaps.
\rightarrow append(1)
                                                                else:
                                                                         ListConstraintsOverlaps.
\rightarrow append(\emptyset)
                                                       else:
                                                                if Y_j - Y_i < i+1:
                                                                         ListConstraintsOverlaps.
\rightarrow append(1)
                                                                else:
                                                                         ListConstraintsOverlaps.
\rightarrow append(\emptyset)
                                              else:
                                                       ListConstraintsOverlaps.append(0)
                  Problem.AddFunction(['sq' + str(i+1), 'sq' + str(j+1)],_
→ListConstraintsOverlaps)
#Problem.Dump('Square.cfn')
                                                                                     (continues on next page)
```

#### 2.12.4 C++ program using libtb2.so

The following code uses the C++ toulbar2 library. Compile toulbar2 with "cmake -DLIBTB2=ON -DPYTB2=ON . ; make" and copy the library in your current directory "cp lib/Linux/libtb2.so ." before compiling "g++ -o square square.cpp -Isrc -Llib/Linux -std=c++11 -O3 -DNDEBUG -DBOOST -DLONGDOUBLE\_PROB - DLONGLONG\_COST -DWCSPFORMATONLY libtb2.so" and running the example (e.g. "./square 15 36").

square.cpp

```
/**
 * Square Packing Problem
 */
// Compile with cmake option -DLIBTB2=ON -DPYTB2=ON to get C++ toulbar2 library lib/
→Linux/libtb2.so
// Then,
// g++ -o square square.cpp -Isrc -Llib/Linux -std=c++11 -O3 -DNDEBUG -DBOOST -
↔ DLONGDOUBLE_PROB -DLONGLONG_COST -DWCSPFORMATONLY libtb2.so
#include "toulbar2lib.hpp"
#include <string.h>
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
int main(int argc, char* argv[])
{
   int N = atoi(argv[1]);
   int S = atoi(argv[2]);
   tb2init(); // must be call before setting specific ToulBar2 options and creating a_
→model
   ToulBar2::verbose = 0; // change to 0 or higher values to see more trace information
   initCosts(); // last check for compatibility issues between ToulBar2 options and.
                                                                            (continues on next page)
```

```
(continued from previous page)
```

```
Generation Generatio Generation 
           Cost top = UNIT_COST;
           WeightedCSPSolver* solver = WeightedCSPSolver::makeWeightedCSPSolver(top);
           for (int i=0; i<N; i++) {
                      solver->getWCSP()->makeEnumeratedVariable(to_string("sq") + to_string(i+1), 0,_
 \hookrightarrow (S-i)*(S-i) - 1);
           }
           for (int i=0; i<N; i++) {
                      for (int j=i+1; j<N; j++) {
                                  vector<Cost> costs((S-i)*(S-i)*(S-j)*(S-j), MIN_COST);
                                             for (int a=0; a<(S-i)*(S-i); a++) {</pre>
                                                         for (int b=0; b<(S-j)*(S-j); b++) {
                                                        costs[a^{*}(S-j)^{*}(S-j)+b] = ((((a^{*}(S-i)) + i + 1 \le (b^{*}(S-j))) || ((b^{*}(S-i))))))
 \rightarrow (S-j)) + j + 1 <= (a%(S-i))) || ((a/(S-i)) + i + 1 <= (b/(S-j))) || ((b/(S-j)) + j + _)
 }
                                  }
                                  solver->getWCSP()->postBinaryConstraint(i, j, costs);
                      }
           }
           solver->getWCSP()->sortConstraints(); // must be done at the end of the modeling
           tb2checkOptions();
           if (solver->solve()) {
                                  vector<Value> sol;
                                  solver->getSolution(sol);
                                             for (int y=0; y<S; y++) {
                                             for (int x=0; x<S; x++) {
                                                        char c = ' ';
                                                        for (int i=0; i<N; i++) {</pre>
                                                                    if (x \ge (sol[i]%(S-i)) \& x < (sol[i]%(S-i)) + i + 1 \& y \ge )
 \hookrightarrow (sol[i]/(S-i)) && y < (sol[i]/(S-i)) + i + 1) {
                                                                               c = 65 + i;
                                                                               break;
                                                                    }
                                                           }
                                                           cout << c;</pre>
                                             }
                                             cout << endl;</pre>
                                  }
           } else {
                                  cout << "No solution found!" << endl;</pre>
           }
           delete solver;
           return 0;
}
```

## 2.13 Square soft packing problem

### 2.13.1 Brief description

The problem is almost identical to the square packing problem with the difference that we now allow overlaps but we want to minimize them.

#### 2.13.2 CFN model

We reuse the *Square packing problem* model except that binary constraints are replaced by cost functions returning the overlapping size or zero if no overlaps.

To calculate an initial upper bound we simply compute the worst case scenario where N squares of size N\*N are all stacked together. The cost of this is N\*\*4, so we will take N\*\*4+1 as the initial upper bound.

#### 2.13.3 Python model

The following code using pytoulbar2 library solves the corresponding cost function network (e.g. "python3 square-soft.py 10 20").

squaresoft.py

```
import sys
from random import randint, seed
seed(123456789)
import pytoulbar2
try:
        N = int(sys.argv[1])
        S = int(sys.argv[2])
        assert N <= S
except:
       print('Two integers need to be given as arguments: N and S')
        exit()
Problem = pytoulbar2.CFN(N^{**4} + 1)
#create a variable for each square
for i in range(N):
        Problem.AddVariable('sq' + str(i+1), ['(' + str(1) + ',' + str(j) + ')' for 1 in_
→range(S-i) for j in range(S-i)])
#binary soft constraints for overlapping squares
for i in range(N):
        for j in range(i+1,N):
                ListConstraintsOverlaps = []
                for a in [S*k+l for k in range(S-i) for l in range(S-i)]:
                        for b in [S*m+n for m in range(S-j) for n in range(S-j)]:
                                 #calculating the coordinates of the squares
                                X_i = a\%S
                                X_j = b\%S
                                Y_i = a//S
                                Y_j = b//S
                                 #calculating if squares are overlapping
```

(continued from previous page) **if** X\_i >= X\_j : if X\_i - X\_j < j+1: **if** Y\_i >= Y\_j: if  $Y_i - Y_j < j+1$ : ListConstraintsOverlaps.  $\Rightarrow$  append(min(j+1-(X\_i - X\_j),i+1)\*min(j+1-(Y\_i - Y\_j),i+1)) else: ListConstraintsOverlaps.  $\rightarrow$  append( $\emptyset$ ) else: **if** Y\_j - Y\_i < i+1: ListConstraintsOverlaps.  $\Rightarrow$  append(min(j+1-(X\_i - X\_j),i+1)\*min(i+1-(Y\_j - Y\_i),j+1)) else: ListConstraintsOverlaps.  $\rightarrow$  append( $\emptyset$ ) else: ListConstraintsOverlaps.append(0) else : **if** X\_j - X\_i < i+1: **if** Y\_i >= Y\_j: **if** Y\_i - Y\_j < j+1: ListConstraintsOverlaps.  $\rightarrow$  append(min(i+1-(X\_j - X\_i), j+1)\*min(j+1-(Y\_i - Y\_j), i+1)) else: ListConstraintsOverlaps.  $\rightarrow$  append( $\emptyset$ ) else: **if** Y\_j - Y\_i < i+1: ListConstraintsOverlaps.  $\rightarrow$  append(min(i+1-(X\_j - X\_i), j+1)\*min(i+1-(Y\_j - Y\_i), j+1)) else: ListConstraintsOverlaps.  $\rightarrow$  append( $\emptyset$ ) else: ListConstraintsOverlaps.append(0) Problem.AddFunction(['sq' + str(i+1), 'sq' + str(j+1)],\_ →ListConstraintsOverlaps) #Problem.Dump('SquareSoft.cfn') Problem.CFN.timer(300) res = Problem.Solve(showSolutions=3) if res: for i in range(S): row = ''for j in range(S): row += ' ' **for** k **in** range(N-1, -1, -1): if  $(res[0][k]%(S-k) \le j \text{ and } j - res[0][k]%(S-k) \le k)_{-}$  $\rightarrow$  and (res[0][k]//(S-k) <= i and i - res[0][k]//(S-k) <= k): row = row[:-1] + chr(65 + k)print(row)

else:
 print('No solution found!')

#### 2.13.4 C++ program using libtb2.so

The following code uses the C++ toulbar2 library. Compile toulbar2 with "cmake -DLIBTB2=ON -DPYTB2=ON . ; make" and copy the library in your current directory "cp lib/Linux/libtb2.so ." before compiling "g++ -o squaresoft squaresoft.cpp -I./src -L./lib/Linux -std=c++11 -O3 -DNDEBUG -DBOOST -DLONGDOUBLE\_PROB -DLONGLONG\_COST -DWCSPFORMATONLY libtb2.so" and running the example (e.g. "./squaresoft 10 20").

squaresoft.cpp

```
/**
* Square Soft Packing Problem
*/
// Compile with cmake option -DLIBTB2=ON -DPYTB2=ON to get C++ toulbar2 library lib/
→Linux/libtb2.so
// Then,
// g++ -o squaresoft squaresoft.cpp -Isrc -Llib/Linux -std=c++11 -O3 -DNDEBUG -DBOOST -
→DLONGDOUBLE_PROB -DLONGLONG_COST -DWCSPFORMATONLY libtb2.so
#include "toulbar2lib.hpp"
#include <string.h>
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
int main(int argc, char* argv[])
{
    int N = atoi(argv[1]);
    int S = atoi(argv[2]);
    tb2init(); // must be call before setting specific ToulBar2 options and creating a_
→model
    ToulBar2::verbose = 0; // change to 0 or higher values to see more trace information
    initCosts(); // last check for compatibility issues between ToulBar2 options and
\hookrightarrowCost data-type
    Cost top = N*(N*(N-1)*(2*N-1))/6 + 1;
    WeightedCSPSolver* solver = WeightedCSPSolver::makeWeightedCSPSolver(top);
    for (int i=0; i < N; i++) {
        solver->getWCSP()->makeEnumeratedVariable(to_string("sq") + to_string(i+1), 0,_
\rightarrow (S-i)*(S-i) - 1);
    }
    for (int i=0; i < N; i++) {
```

```
for (int j=i+1; j < N; j++) {
                                   vector<Cost> costs((S-i)*(S-i)*(S-j)*(S-j), MIN_COST);
                                               for (int a=0; a < (S-i)*(S-i); a++) {</pre>
                                                          for (int b=0; b < (S-j)*(S-j); b++) {
                                                          costs[a^{*}(S-j)^{*}(S-j)+b] = ((((a^{(S-i))} + i + 1 \le (b^{(S-j)}))) || ((b^{(S-i)}))) || ((b^{(S-i)})) |
 \rightarrow (S-j)) + j + 1 <= (a%(S-i))) || ((a/(S-i)) + i + 1 <= (b/(S-j))) || ((b/(S-j)) + j + _)
 \rightarrow 1 \le (a/(S-i)))?MIN_COST:(min((a%(S-i)) + i + 1 - (b%(S-j)), (b%(S-j)) + j + 1 - (a)))
 \rightarrow (S-i)) * min((a/(S-i)) + i + 1 - (b/(S-j)), (b/(S-j)) + j + 1 - (a/(S-i))));
                                               }
                                   }
                                   solver->getWCSP()->postBinaryConstraint(i, j, costs);
                       }
           }
           solver->getWCSP()->sortConstraints(); // must be done at the end of the modeling
           tb2checkOptions();
           if (solver->solve()) {
                                   vector<Value> sol;
                                   solver->getSolution(sol);
                                               for (int y=0; y < S; y++) {
                                               for (int x=0; x < S; x++) {
                                                          char c = ' ';
                                                          for (int i=N-1; i >= 0; i--) {
                                                                      if (x >= (sol[i]%(S-i)) && x < (sol[i]%(S-i) ) + i + 1 && y >=_
 ⇔(sol[i]/(S-i)) && y < (sol[i]/(S-i)) + i + 1) {
                                                                                  if (c != ' ') {
                                                                                              c = 97 + i;
                                                                                  } else {
                                                                                              c = 65+i;
                                                                                  }
                                                                      }
                                                              }
                                                              cout << c;</pre>
                                               }
                                               cout << endl;</pre>
                                   }
           } else {
                                   cout << "No solution found!" << endl;</pre>
           }
           delete solver;
           return 0;
}
```

## 2.14 Golomb ruler problem

#### 2.14.1 Brief description

A golomb ruler of order N is a set of integer marks 0=a1<a2<a3<a4<...<aN such that each difference between two ak's is unique.

For example, this is a golomb ruler:



We can see that all differences are unique, rather than in this other ruler where 0-3 and 3-6 are both equal to 3.



The size of a golomb ruler is equal to aN, the greatest number of the ruler. The goal is to find the smallest golomb ruler given N.

#### 2.14.2 CFN model

We create N variables, one for each integer mark ak. Because we can not create an AllDifferent constraint with differences of variables directly, we also create a variable for each difference and create hard ternary constraints in order to force them be equal to the difference. Because we do not use an absolute value when creating the hard constraints, it forces the assignment of ak's variables to follow an increasing order.

Then we create an AllDifferent constraint on all the difference variables and one unary cost function on the last aN variable in order to minimize the size of the ruler. In order to break symmetries, we set the first mark to be zero.

#### 2.14.3 Python model

The following code using pytoulbar2 library solves the golomb ruler problem with the first argument being the number of marks N (e.g. "python3 golomb.py 8").

golomb.py

```
import sys
import pytoulbar2
N = int(sys.argv[1])
top = N^{**2} + 1
Problem = pytoulbar2.CFN(top)
#create a variable for each mark
for i in range(N):
   Problem.AddVariable('X' + str(i), range(N**2))
#ternary constraints to link new variables of difference with the original variables
for i in range(N):
    for j in range(i+1, N):
        Problem.AddVariable('X' + str(j) + '-X' + str(i), range(N**2))
        Constraint = []
        for k in range(N**2):
            for 1 in range(N**2):
                for m in range(N**2):
```

```
if 1-k == m:
                        Constraint.append(0)
                    else:
                         Constraint.append(top)
        Problem.AddFunction(['X' + str(i), 'X' + str(j), 'X' + str(j) + '-X' + str(i)], \_
→Constraint)
Problem.AddAllDifferent(['X' + str(j) + '-X' + str(i) for i in range(N) for j in_
\rightarrowrange(i+1,N)])
Problem.AddFunction(['X' + str(N-1)], range(N**2))
#fix the first mark to be zero
Problem.AddFunction(['X0'], [0] + [top] * (N**2 - 1))
#Problem.Dump('golomb.cfn')
Problem.CFN.timer(300)
res = Problem.Solve(showSolutions=3)
if res:
    ruler = '0'
    for i in range(1,N):
        ruler += ' '*(res[0][i]-res[0][i-1]-1) + str(res[0][i])
    print('Golomb ruler of size:',int(res[1]))
    print(ruler)
```

## 2.15 Board coloration problem

#### 2.15.1 Brief description

Given a rectangular board with dimension n\*m, the goal is to color the cells such that any inner rectangle included inside the board doesn't have all its corners colored with the same color. The goal is to minimize the number of colors used.

For example, this is not a valid solution of the 3\*4 problem, because the red and blue rectangles have both their 4 corners having the same color:



On the contrary, the following coloration is a valid solution of the 3\*4 problem because every inner rectangle inside

the board does not have a unique color for its corners:



#### 2.15.2 CFN basic model

We create n\*m variables, one for each square of the board, with domain size equal to n\*m representing all the possible colors. We also create one variable for the number of colors.

We create hard quaternary constraints for every rectangle inside the board with a cost equal to 0 if the 4 variables have different values and a forbidden cost if not.

We then create hard binary constraints between the variable of the number of colors for each cell to fix the variable for the number of colors as an upper bound.

Then we create a soft constraint on the number of colors to minimize it.

#### 2.15.3 Python model

The following code using pytoulbar2 library solves the board coloration problem with the first two arguments being the dimensions n and m of the board (e.g. "python3 boardcoloration.py 3 4").

boardcoloration.py

```
import sys
from random import randint, seed
seed(123456789)
import pytoulbar2
try:
    n = int(sys.argv[1])
    m = int(sys.argv[2])
except:
    print('Two integers need to be in arguments: number of rows n, number of columns m')
    exit()
top = n*m + 1
Problem = pytoulbar2.CFN(top)
#create a variable for each cell
for i in range(n):
    for j in range(m):
```

```
(continued from previous page)
        Problem.AddVariable('sq_' + str(i) + '_' + str(j), range(n*m))
#create a variable for the maximum of colors
Problem.AddVariable('max', range(n*m))
#quaterny hard constraints for rectangle with same color angles (encoding with forbidden...
\rightarrow tuples)
ConstraintTuples = []
ConstraintCosts = []
for k in range(n*m):
    #if they are all the same color
   ConstraintTuples.append([k, k, k, k])
   ConstraintCosts.append(top)
#for each cell on the chessboard
for i1 in range(n):
    for i2 in range(m):
        #for every cell on the chessboard that could form a valid rectangle with the
→ first cell as up left corner and this cell as down right corner
        for j1 in range(i1+1, n):
            for j2 in range(i2+1, m):
                # add a compact function with zero default cost and only forbidden tuples
                Problem.AddCompactFunction(['sq_' + str(i1) + '_' + str(i2), 'sq_' +_

str(i1) + '_' + str(j2), 'sq_' + str(j1) + '_' + str(i2), 'sq_' + str(j1) + '_' +

str(j2)], 0, ConstraintTuples, ConstraintCosts)

#binary hard constraints to fix the variable max as an upper bound
Constraint = []
for k in range(n*m):
    for 1 in range(n*m):
        if k>1:
            #if the color of the square is more than the number of the max
            Constraint.append(top)
        else:
            Constraint.append(0)
for i in range(n):
    for j in range(m):
        Problem.AddFunction(['sq_' + str(i) + '_' + str(j), 'max'], Constraint)
#minimize the number of colors
Problem.AddFunction(['max'], range(n*m))
#symmetry breaking on colors
for i in range(n):
    for j in range(m):
        Constraint = []
        for k in range(n*m):
            if k > i^{*}m+j:
                Constraint.append(top)
            else:
                Constraint.append(♥)
        Problem.AddFunction(['sq_' + str(i) + '_' + str(j)], Constraint)
```

```
#Problem.Dump('boardcoloration.cfn')
Problem.CFN.timer(300)
res = Problem.Solve(showSolutions = 3)
if res:
    for i in range(n):
        row = []
        for j in range(m):
            row.append(res[0][m*i+j])
        print(row)
else:
    print('No solution found!')
```

## 2.16 Learning to play the Sudoku

#### 2.16.1 Available

• Presentation



## 2.17 Learning car configuration preferences

#### 2.17.1 Brief description

Renault car configuration system: learning user preferences.

#### 2.17.2 Available

• Presentation



• Data GitHub code

## 2.18 Visual Sudoku Tutorial

#### 2.18.1 Brief description

A simple case mixing Deep Learning and Graphical models.

## 2.18.2 Available

• You can run it directly from your browser as a Jupyter Notebook



## 2.19 Visual Sudoku Application

#### 2.19.1 Brief description

An automatic Sudoku puzzle solver using OpenCV, Deep Learning, and Optical Character Recognition (OCR).

#### 2.19.2 Available

#### Software

Software adapted by Simon de Givry (@ INRAE, 2022) in order to use toulbar2 solver, from a tutorial by Adrian

Rosebrock (@ PyImageSearch, 2022) : GitHub code

#### As an APK

Based on this software, a 'Visual Sudoku' application for Android has been developed to be used from a smartphone.

See the *detailed presentation* (description, source, download...).

The application allows to capture a grid from its own camera ('*CAMERA' menu*) or to select a grid among the smartphone existing files ('*FILE' menu*), for example files coming from 'DCIM', in .jpg or .png formats. The grid image must have been captured in portrait orientation. Once the grid has been chosen, the '*Solve'* button allows to get the solution.



Fig.1

- Fig.1 : Screen of main menu
- Fig.2 : Screen of the grid to be solved
- Fig.3 : Screen of the solution (in yellow) found by the solver

Examples of some input grids and their solved grids

#### As a Web service

The software is available as a web service. The **visual sudoku web service**, hosted by the ws web services (based on HTTP protocol), can be called by many ways : from a **browser** (like above), from any softwares written in a language supporting HTTP protocol (**Python**, **R**, **C++**, **Java**, **Php**...), from command line tools (**cURL**...)...

• Calling the visual sudoku web service from a browser :



• Example of calling the visual sudoku web service from a terminal by cURL :

Commands (replace mygridfilename.jpg by your own image file name) :

```
curl --output mysolutionfilename.jpg -F 'file=@mygridfilename.jpg' -F 'keep=40' -F

→'border=15' http://147.100.179.250/api/tool/vsudoku
```

• The 'Visual Sudoku' APK calls the visual sudoku web service.

## 2.20 Visual Sudoku App for Android

#### 2.20.1 A visual sudoku solver based on cost function networks

This application solves the sudoku problem from a smartphone by reading the grid using its camera. The cost function network solver toulbar2 is used to deal with the uncertainty on the digit recognition produced by the neural network. This uncertainty, combined with the sudoku logical rules, makes it possible to correct perceptual errors. It is particularly useful in the case of hand-written digits or poor image quality. It is also possible to solve a partially filled-in grid with printed and hand-written digits. The solver will always suggest a valid solution that best adapts to the retrieved digit information. It will naturally detect (a small number of) errors in a partially filled-in grid and could be used later as a diagnosis tool (future work). This software demonstration emphasizes the tight relation between constraint programming, computer vision, and deep learning.

We used the open-source C++ solver toulbar2 in order to find the maximum a posteriori solution of a constrained probabilistic graphical model. With its dedicated numerical (soft) local consistency bounds, toulbar2 outperforms traditional CP solvers on this problem. Grid perception and cell extraction are performed by the computer vision library OpenCV. Digit recognition is done by **Keras** and **TensorFlow**. The current android application is written in Python using the Kivy framework. It is inspired from a tutorial by Adrian Rosebrock. It uses the ws RESTful web services in order to run the solver.

See also : Visual Sudoku Application.

#### 2.20.2 Source Code

GitHub code

#### 2.20.3 Download and Install

To install the 'Visual Sudoku' application on smartphone :

1) Download the visualsudoku-release.apk APK file from Github repository :



#### https://github.com/toulbar2/visualsudoku/releases/latest

- 2) Click on the downloaded **visualsudoku-release.apk** APK file to ask for **installation** (*you have to accept to 'install anyway' from unknown developer*).
- 3) In your parameter settings for the app, give permissions to the 'Visual Sudoku' application (smartphone menu 'Parameters' > 'Applications' > 'Visual Sudoku') : allow camera (required to capture grids), files and multimedia contents (required to save images as files). Re-run the app.

Warnings :

- The application may fail at first start and you may have to launch it twice.
- While setting up successfully, the application should have created itself the required 'VisualSudoku' folder (under the smartphone 'Internal storage' folder) but if not, you will have to create it by yourself manually.
- Since the application calls a web service, an internet connection is required.

#### 2.20.4 Description

The 'SETTINGS' menu allows to save grids or solutions as image files ('savinginputfile', 'savingoutputfile' parameters) and to access to some 'expert' parameters in order to enhance the resolution process ('keep', 'border', 'time' parameters).

The application allows to capture a grid from its own camera (*'CAMERA' menu*) or to select a grid among the smartphone existing files (*'FILE' menu*), for example files coming from 'DCIM', in .jpg or .png formats. The grid image must have been captured in portrait orientation. Once the grid has been chosen, the *'Solve'* button allows to get the solution.

Fig.2



8|7

Fig.1

Fig.3



- Fig.1 : Screen of main menu
- Fig.2 : Screen of the grid to be solved
- Fig.3 : Screen of the solution (in yellow) found by the solver

Examples of some input grids and their solved grids

## 2.21 A sudoku code

#### 2.21.1 Brief description

A Sudoku code returning a sudoku partial grid (sudoku problem) and the corresponding completed grid (sudoku solution), such as partial and completed grids.

The verbose version, that further gives a detailed description of what the program does, could be useful as tutorial example. Example : partial and completed grids with explanations.

#### 2.21.2 Available

Available as a web service.

You can run the software directly from your browser as a web service :

Grids information is returned into the output stream. The **returned\_type** parameter of the web service allows to choose how to receive it :

- *returned\_type=stdout.txt* : to get the output stream as a .txt file.
- *returned\_type=run.zip* : to get the .zip run folder containing the output stream \_\_WS\_\_stdout.txt (+ *the error stream \_\_WS\_\_stderr.txt that may be useful to investigate*).

Web service to get one sudoku grids (both partial and completed) :



Web service to further get a detailed description of what the program does (verbose version) :



#### Note

The **sudoku web services**, hosted by the ws web services (based on HTTP protocol), can be called by many other ways : from a **browser** (like above), from any softwares written in a language supporting HTTP protocol (**Python**, **R**, **C++**, **Java**, **Php**...), from command line tools (**cURL**...)...

#### Example of calling the sudoku web services from a terminal by cURL :

• Commands (replace indice value by any value in 1...17999):

```
curl --output mygrids.txt -F 'indice=778' -F 'returned_type=stdout.txt' http://147.

→100.179.250/api/tool/sudoku

curl --output myrun.zip -F 'indice=778' -F 'returned_type=run.zip' http://147.100.

→179.250/api/tool/sudoku

# verbose version

curl --output mygrids_details.txt -F 'indice=778' -F 'returned_type=stdout.txt'__

(continues on next page)
```

- Responses corresponding with the requests above :
  - mygrids.txt
  - \_\_WS\_\_stdout.txt into myrun.zip has the same content as mygrids.txt
  - mygrids\_details.txt (\_\_WS\_\_stdout.txt into myrun\_details.zip has the same content)
  - \_WS\_stdout.txt into myrun\_details.zip has the same content as mygrids\_details.txt